

IP Next Generation (IPv6)

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- Limitations of current Internet Protocol (IP)
- How many addresses do we need?
- IPv6 Addressing
- IPv6 header format
- IPv6 features: routing flexibility, plug-n-play, multicast support, flows

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IP Addresses

- ❑ **Example:** 164.107.134.5
= 1010 0100 : 0110 1011 : 1000 0110 : 0000 0101
= A4:6B:86:05 (32 bits)
- ❑ Maximum number of address = $2^{32} = 4$ Billion
- ❑ Class A Networks: 15 Million nodes
- ❑ Class B Networks: 64,000 nodes or less
- ❑ Class C Networks: 250 nodes or less

IP Address Format

- ❑ Three all-zero network numbers are reserved
- ❑ 127 Class A + 16,381 Class B + 2,097,151 Class C networks = 2,113,659 networks total
- ❑ Class B is most popular.
- ❑ 20% of Class B were assigned by 7/90 and doubling every 14 months \Rightarrow Will exhaust by 3/94
- ❑ Question: Estimate how big will you become?
Answer: More than 256!
Class C is too small. Class B is just right.

How Many Addresses?

- ❑ 10 Billion people by 2020
- ❑ Each person has more than one computer
- ❑ Assuming 100 computers per person $\Rightarrow 10^{12}$ computers
- ❑ More addresses may be required since
 - ❑ Multiple interfaces per node
 - ❑ Multiple addresses per interface
- ❑ Some believe 2^6 to 2^8 addresses per host
- ❑ Safety margin $\Rightarrow 10^{15}$ addresses
- ❑ IPng Requirements $\Rightarrow 10^{12}$ end systems and 10^9 networks. Desirable 10^{12} to 10^{15} networks

Address Size

- ❑ H Ratio = $\log_{10}(\# \text{ of objects})/\text{available bits}$
- ❑ 2^n objects with n bits: H-Ratio = $\log_{10}2 = 0.30103$
- ❑ French telephone moved from 8 to 9 digits at 10^7 households $\Rightarrow H = 0.26$ (~3.3 bits/digit)
- ❑ US telephone expanded area codes with 10^8 subscribers $\Rightarrow H = 0.24$
- ❑ Physics/space science net stopped at 15000 nodes using 16-bit addresses $\Rightarrow H = 0.26$
- ❑ 3 Million Internet hosts currently using 32-bit addresses $\Rightarrow H = 0.20 \Rightarrow$ A few more years to go

IPv6 Addresses

- ❑ 128-bit long. Fixed size
- ❑ $2^{128} = 3.4 \times 10^{38}$ addresses
 - ⇒ 665×10^{21} addresses per sq. m of earth surface
- ❑ If assigned at the rate of $10^6/\mu\text{s}$, it would take 20 years
- ❑ Expected to support 8×10^{17} to 2×10^{33} addresses
 - $8 \times 10^{17} \Rightarrow 1,564$ address per sq. m
- ❑ Allows multiple interfaces per host.
- ❑ Allows multiple addresses per interface
- ❑ Allows unicast, multicast, anycast
- ❑ Allows provider based, site-local, link-local
- ❑ 85% of the space is unassigned

Colon-Hex Notation

- ❑ **Dot-Decimal:** 127.23.45.88
- ❑ **Colon-Hex:**
FEDC:0000:0000:0000:3243:0000:0000:ABCD
 - ❑ Can skip leading zeros of each word
 - ❑ Can skip one sequence of zero words, e.g.,
FEDC::<3243:0000:0000:ABCD *or*
::3243:0000:0000:ABCD
 - ❑ Can leave the last 32 bits in dot-decimal, e.g.,
::127.23.45.88
 - ❑ Can specify a prefix by /length, e.g.,
2345:BA23:7::/40

Header

□ IPv6:

Version	Class	Flow Label
Payload Length		Next Header Hop Limit
Source Address		
Destination Address		

□ IPv4:

Version	IHL	Type of Service	Total Length
Identification		Flags	Fragment Offset
Time to Live	Protocol	Header Checksum	
Source Address			
Destination Address			
Options			Padding

IPv6 vs IPv4

- IPv6 only twice the size of IPv4 header
- Only version number has the same position and meaning as in IPv4
- **Removed:**
 - Header length, fragmentation fields (identification, flags, fragment offset), header checksum
- **Replaced:**
 - Datagram length by payload length
 - Protocol type by next header
 - Time to live by hop limit
 - Type of service by “class” octet
- **Added:** flow label
- All fixed size fields.

IPv6 vs IPv4

- ❑ No optional fields. Replaced by extension headers.
 - ❑ Idea: avoid unnecessary processing by intermediate routers while not sacrificing the flexibility possible due to options
 - ❑ Next Header = 6 (TCP), 17 (UDP), etc

Extension Headers



- ❑ Most extension headers are examined only at destination
- ❑ Routing: Loose or tight source routing
- ❑ Fragmentation: All IPv6 routers can carry 536 Byte payload
- ❑ Authentication
- ❑ Hop-by-Hop Options
- ❑ Destination Options:

Extension Header (Cont)

- ❑ Only Base Header:

Base Header Next = TCP	TCP Segment
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- ❑ Only Base Header and One Extension Header:

Base Header Next = TCP	Route Header Next = TCP	TCP Segment
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- ❑ Only Base Header and Two Extension Headers:

Base Header Next = TCP	Route Header Next = Auth	Auth Header Next = TCP	TCP Segment
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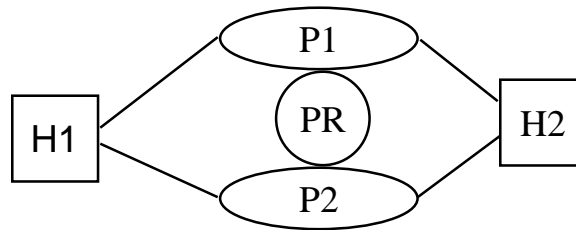
Routing Header

Next Header	Hdr Ext Len	Routing Type	Sgmts left
Reserved			
Address 1			
Address 2			
Address n			

- ❑ Type = 0 ⇒ Current source routing
- ❑ Router will look at RH if its address is in the destination field
- ❑ New Functionality: Provider selection, Host mobility, Auto-readdressing (route to new address)

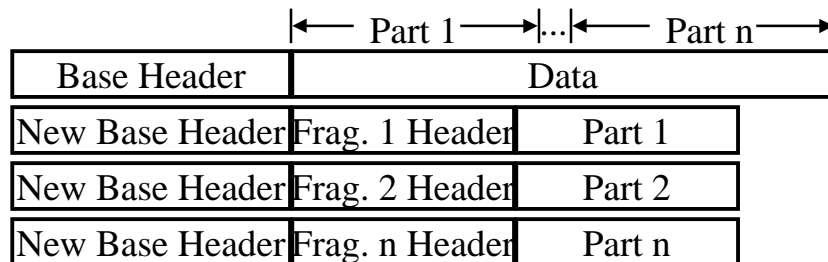
Provider Selection

- ❑ Possible using routing extension header
- ❑ Source specified intermediate systems
- ❑ No preference: H1, H2
- ❑ P1 Preferred: H1, P1, H2
- ❑ H1 becomes Mobile: H1, PR, P1, H2



Fragmentation

- ❑ Routers cannot fragment. Only source hosts can.
 - ⇒ Need path MTU discovery or tunneling
- ❑ Fragmentation requires an extension header
- ❑ Payload is divided into pieces
- ❑ A new base header is created for each fragment



IPv6 addressing and routing

- ❑ Aggregatable Global Unicast Addresses
- ❑ Link-local and Site-local addresses
- ❑ Multicast and Anycast support
- ❑ Provider-based inter-domain routing & IDR

Initial IPv6 Prefix Allocation

Allocation	Prefix	Allocation	Prefix
Reserved	0000 0000	Unassigned	101
Unassigned	0000 0001	Unassigned	110
NSAP	0000 001	Unassigned	1110
IPX	0000 010	Unassigned	1111 0
Unassigned	0000 011	Unassigned	1111 10
Unassigned	0000 1	Unassigned	1111 110
Unassigned	0001	Unassigned	1111 1110
Unassigned	001	Unassigned	1111 1110 0
Provider-based*	010	Link-Local	1111 1110 10
Unassigned	011	Site-Local	1111 1110 11
Geographic	100	Multicast	1111 1111

*Has been renamed as “Aggregatable global unicast”

Aggregatable Global Unicast Addresses

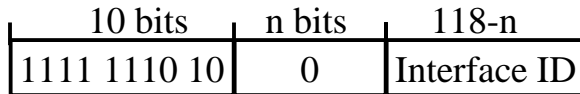
- ❑ Address allocation: “provider-based” plan
- ❑ Format: TLA + NLA + SLA + 64-bit interface ID
- ❑ TLA = “Top level aggregator.” (13 bits)
 - ❑ Ranges of TLA values allocated to various registries
 - ❑ For “backbone” providers or “exchange points”
- ❑ NLA = “Next Level Aggregator” (32 bits)
 - ❑ Second tier provider and a subscriber
 - ❑ More levels of hierarchy possible within NLA
- ❑ SLA = “Site level aggregator” = 16 bits for link

Aggr. Global Unicast Addr

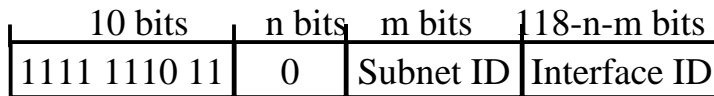
- ❑ Renumbering after change of provider => change the TLA and NLA. But have same SLA & I/f ID
- ❑ Interface ID = 64 bits
 - ❑ Will be based on IEEE EUI-64 format
 - ❑ An extension of the IEEE 802 (48 bit) format.
 - ❑ Possible to derive the IEEE EUI-64 equivalent of current IEEE 802 addresses
 - ❑ Along with neighbor discovery procedures, obviates need for ARP.

Local-Use Addresses

- ❑ Link Local: Not forwarded outside the link,
FE:80::xxx

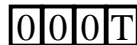
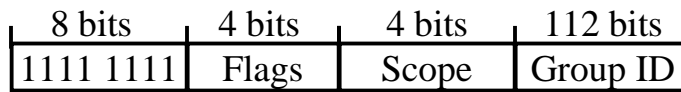


- ❑ Site Local: Not forwarded outside the site,
FE:C0::xxx



- ❑ Provides plug and play

Multicast Addresses



- ❑ All routers recognize this format => all routers can route multicast packets. Also IGMP part of ICMPv6 => required.
- ❑ Flags: T = 0 => Permanent (well-known) multicast address, 1 => Transient
- ❑ Scope: 1 Node-local, 2 Link-local, 5 Site-local, 8 Organization-local, E Global => routers reqd to honor this.
- ❑ Predefined: 1 => All nodes, 2 => Routers, 1:0 => DHCP servers

Multicast & Anycast

- ❑ Example: 43 ⇒ NTP Servers
 - ❑ FF01::43 ⇒ All NTP servers on this node
 - ❑ FF02::43 ⇒ All NTP servers on this link
 - ❑ FF05::43 ⇒ All NTP servers in this site
 - ❑ FF08::43 ⇒ All NTP servers in this org.
 - ❑ FF0F::43 ⇒ All NTP servers in the Internet
- ❑ Structure of Group ID:
 - ❑ First 80 bits = zero (to avoid risk of group collision, because IP multicast mapping uses only 32 bits)

Address Autoconfiguration

- ❑ Allows plug and play
- ❑ BOOTP and DHCP are used in IPv4
- ❑ DHCPng will be used with IPv6
- ❑ Two Methods: Stateless and Stateful
- ❑ Stateless:
 - ❑ A system uses link-local address as source and multicasts to "All routers on this link"
 - ❑ Router replies and provides all the needed prefix info

Address Autoconfiguration

- ❑ All prefixes have a associated lifetime
- ❑ System can use link-local address permanently if no router
- ❑ Stateful:
 - ❑ Problem w stateless: Anyone can connect
 - ❑ Routers ask the new system to go DHCP server (by setting managed configuration bit)
 - ❑ System multicasts to "All DHCP servers"
 - ❑ DHCP server assigns an address

Neighbor Discovery

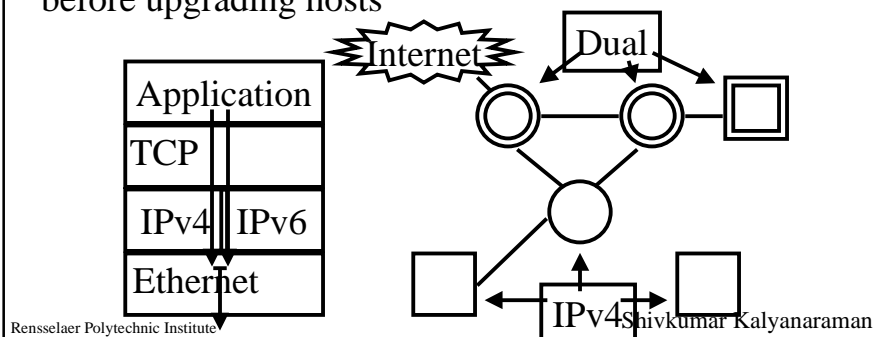
- ❑ Media addresses allowed to be upto 128 bits long
- ❑ Part of ICMPv6 functionality
- ❑ Subsumes ARP, Router discovery.
- ❑ Source maintains several caches: destination cache, neighbor cache, prefix cache, router cache
- ❑ Multicast solicitation for neighbor media address if unavailable in neighbor cache
- ❑ Neighbor advertisement message sent to soliciting station.
- ❑ Redirects also part of ICMPv6

Real-time support

- ❑ Flow label and the “class” octet field
- ❑ Flow = sequence of packets from a single source to a particular (unicast/multicast) destinations requiring special handling by intermediate routers
- ❑ Applications becoming adaptive
 - ❑ Even adaptive voice available for IP telephony
- ❑ Hierarchical transmissions:
 - ❑ Can cause congestion {Steve McCanne, SIGCOMM’96}
=>“priority” renamed as “class”
 - ❑ “Class” field currently being worked upon by differentiated services group

Transition Mechanisms

- ❑ Dual-IP Hosts, Routers, Name servers
- ❑ Tunneling IPv6 over IPv4
- ❑ Nodes can be partially upgraded to IPv6
- ❑ It is better (though not required) to upgrade routers before upgrading hosts



Application Issues

- ❑ Most application protocols will have to be upgraded: FTP, SMTP, Telnet, Rlogin
- ❑ 27 of 51 Full Internet standards, 6 of 20 draft standards, 25 of 130 proposed standards will be revised for IPv6
- ❑ No checksum \Rightarrow checksum at upper layer is mandatory, even in UDP
- ❑ non-IETF standards: X-Open, Kerberos, ... will be updated
- ❑ Should be able to request and receive new DNS records

Summary



- ❑ IPv6 uses 128-bit addresses
- ❑ Allows provider-based, site-local, link-local, multicast, anycast addresses
- ❑ Fixed header size. Extension headers instead of options for provider selection, security etc
- ❑ Allows auto-configuration
- ❑ Dual-IP router and host implementations for transition