

Better-than-best-effort: QoS, Int-serv, Diff-serv, RSVP, RTP

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- QoS building blocks
- ATM QoS architecture
- Why better-than-best-effort Internet ?
- Support for multimedia apps: RTP, H.323, Integrated Services(int-serv), RSVP.
- Scalable differentiated services for ISPs: diff-serv
- Missing pieces: QoS routing, traffic engineering, policy management, pricing models

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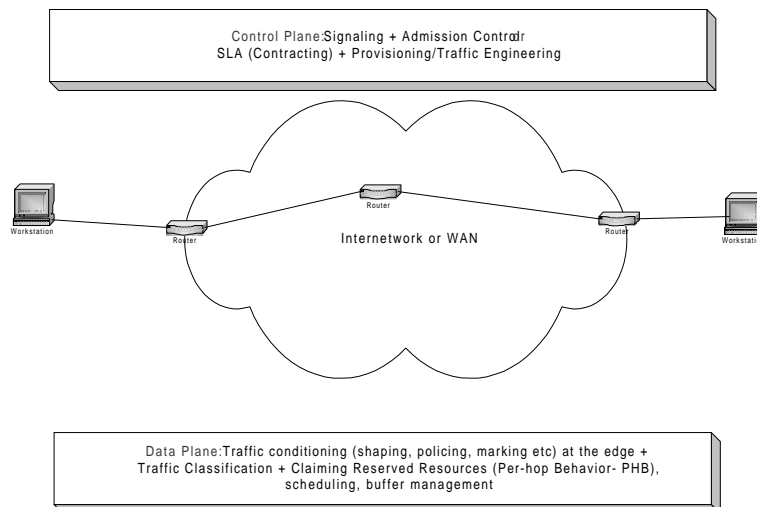
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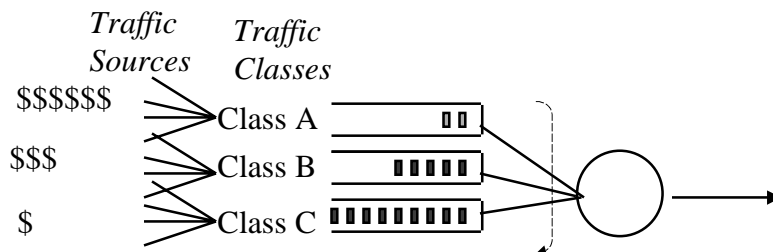
QoS building blocks

- ❑ QoS => set aside resources for premium services
- ❑ QoS components:
 - ❑ a) What kind of premium services ? (*Service/SLA design*)
 - ❑ b) How much resources? (*admission control/provisioning*)
 - ❑ c) How to ensure network resource utilization, do load balancing, flexibly manage traffic aggregates and paths ? (*QoS routing, traffic engineering*)
 - ❑ d) How to actually set aside these resources in a distributed manner ? (*signaling, provisioning, policy*)
 - ❑ e) How to deliver the service when the traffic actually comes in ? (*traffic shaping, classification, scheduling*)
 - ❑ f) How to monitor quality, account and price these services? (*Network management, Accounting, Billing, Pricing*)

QoS big picture: Control/Data planes



Eg. Mechanisms: Queuing/Scheduling



- Use a few bits to indicate which queue (class) a packet goes into (also branded as CoS)
- High \$\$ users get into high priority queues, which are in turn less populated => lower delay and near-zero likelihood of packet drop

Eg. Mechanisms (contd): priority drop

Drop In and out-of-profile packets



Drop only out-of-profile packets

- Enhance buffer management to preferentially drop red packets when a low threshold is crossed

ATM QoS framework

- ❑ **Services:** CBR, rt-VBR, nrt-VBR, ABR, UBR
- ❑ **QoS Routing and Signaling:**
 - ❑ PNNI, ATM signaling with VCs/VPs
- ❑ **Traffic management:**
 - ❑ QoS parameter design, traffic conditioners, feedback control
 - ❑ Standard end system and switch behavior for each of the services
- ❑ **Critique:** No support for qualitative, provider-defined services, limited pt-to-mpt support

ATM Traffic Classes

- ❑ **CBR, VBR** for voice, video: “higher priority”
- ❑ **ABR, GFR, UBR** for data: uses “left over capacity”
- ❑ **ABR properties:** low latency, high throughput, fairness among contending sources, and low cell loss.
- ❑ **UBR properties:** No guarantees. Happy-go-lucky.
- ❑ **GFR properties:**
 - ❑ Minimum rate provided through simple signaling and buffer management.
 - ❑ Intermediate to ABR and UBR - similar to frame relay.

Internet real-time support model

- ❑ Initially assume that the net offers no real-time support and engineer transport protocols (RTP) and middleware which can enable adaptive real-time applications
- ❑ On the longer term, build QoS mechanisms: control-plane and data-plane
- ❑ Flexibility to leverage the Internet connectionless model, allow for future multicast capability, accommodate ISP's desire to "provision/engineer" networks, and design their own services

RTP

- ❑ RTP is the standard protocol for the transport of real-time data, including audio and video.
- ❑ RTP follows the application level framing (ALF) philosophy.
 - ❑ RTP specifies common app functions.
 - ❑ It is intended to be tailored through modifications and/or additions to the headers (spec'd in companion docs)
- ❑ RTP consists of a data and a control part. The latter is called RTCP.
- ❑ The data part of RTP is a thin protocol.

RTCP

- ❑ RTCP provides support for real-time conferencing of groups of any size within an internet.
 - ❑ Eg: source identification and support for gateways like audio and video bridges as well as multicast-to-unicast translators.
 - ❑ It offers quality-of-service feedback from receivers to the multicast group & synchronization support for media streams.

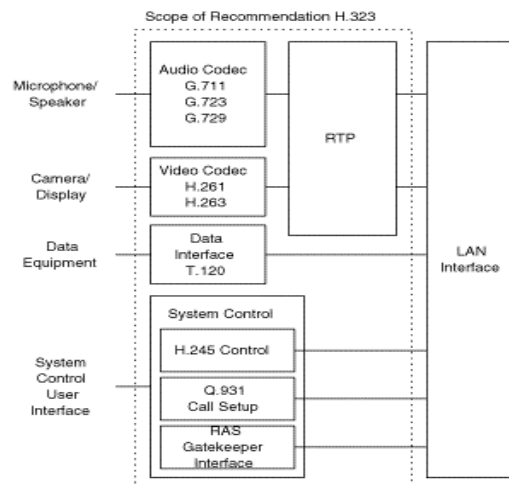
RTP (contd)

- ❑ RTP services: payload type identification, sequence numbering, timestamping, delivery monitoring, & optional mixing/translation. UDP for multiplexing and checksum services
- ❑ RTP does not provide: mechanisms to ensure quality-of-service, guarantee delivery or prevent out-of-order delivery or loss.
 - ❑ RTP sequence numbers allow receiver to reconstruct the sender's packet sequence, or to determine the proper location of a packet, eg, in video decoding, without necessarily decoding packets in sequence.

H.323

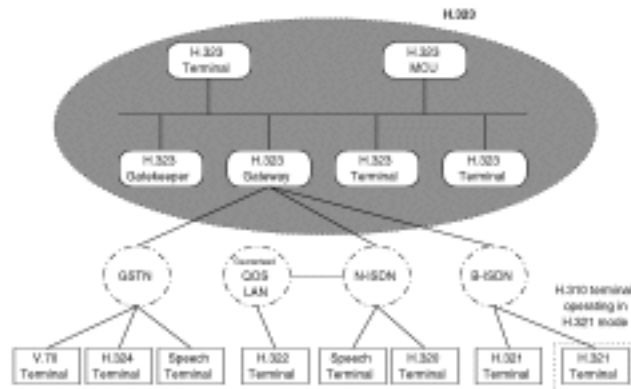
- ❑ H.323 is an ITU standard for multimedia communications over best-effort LANs.
- ❑ Part of larger set of standards (H.32X) for videoconferencing over data networks.
- ❑ H.323 includes both stand-alone devices and embedded personal computer technology as well as point-to-point and multipoint conferences.
- ❑ H.323 addresses call control, multimedia management, and bandwidth management as well as interfaces between LANs and other networks.

H.323 Architecture



H.323 (contd)

- ❑ Terminals, Gateways, Gatekeepers, and Multipoint Control Units (MCUs)



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H.323 (contd)

- ❑ Terminals: All terminals must support voice; video and data are optional.
- ❑ Gateway: an optional element which provides translation functions between H.323 conferencing endpoints (esp for ISDN, PSTN)
- ❑ Gatekeeper: most important component which provides call control services
- ❑ Multipoint Control Unit (MCU): supports conferences between three or more endpoints. Consists of a Multipoint Controller (MC) and Multipoint Processors (MP).

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Integrated Services (int-serv)

- Supplement Internet Architecture with:
 - Services: guaranteed delay, controlled load
 - New signaling protocol: RSVP + admission control
 - *Shaping* at edge nodes combines with packet *classification* and *scheduling/buffer management* at routers to provide local delay and bandwidth guarantees.
 - Specs for parameters (flow-spec), classification (filter-spec)
- **Critique**: non-scalable, no control over routing vagaries, no feedback support

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RSVP

- A *signaling* protocol: creates and maintains distributed reservation state
- *Multicast* trees setup by routing protocols, not RSVP (unlike ATM signaling)
- *Receiver-initiated*: scales for multicast
- *Soft-state*: time out unless refreshed: robust.
- Latest paths discovered through "PATH" messages and used by RESV mesgs.
- *Flowspec*: specifies resource to be reserved
- *Filterspec*: specifies how to classify packets
- Reservation styles: "wildcard", "fixed-filter", and "dynamic-filter".

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Diff-serv motivations

- ❑ #1. *Economics* of ISPs (access and transit providers) dictates *need for service differentiation*
 - ❑ IP provides just a best effort service
 - ❑ TOS is used in a non-standard way, and could be redefined to be more useful
 - ❑ Work done in pricing aspects of SLAs did not fit into IP because of a lack of header bits
 - ❑ ISPs, not IETF, should define services
 - ❑ Some services could be end-to-end, but here IETF would standardize only building blocks

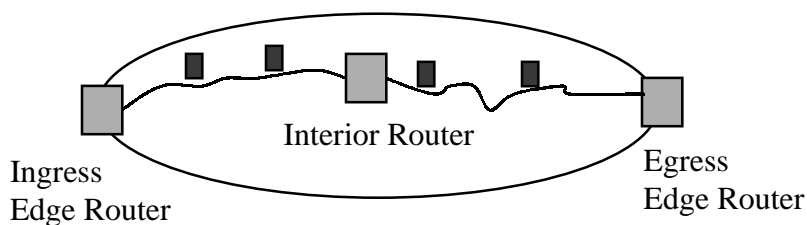
Diff-serv motivations (contd)

- ❑ #2. Diffserv is considered to be *crucial building block* to provide performance assurances *in IP-based VPNs*.
 - ❑ Other pieces: IPSEC (security & tunneling), L2TP (remote-access tunneling), and RSVP (QoS signaling)
- ❑ #3. *Int-serv/RSVP does not scale*
 - ❑ Diff-serv uses a limited set of “behavior aggregates (BA)”
 - ❑ Diffserv creates a separation between edge and core routers.
 - ❑ Move per-flow (possibly non-scalable) data path functions (or MF-classification) to edges.
 - ❑ Edge handles policy, contracting and billing.
 - ❑ Interiors may participate in signaling

Diff-serv motivations (contd)

- *Diff-serv must work with IPv4.*
 - Costs: incompatibility...
 - Redefining TOS octet.
 - Compatibility w/ RFC 791 (IP precedence)
 - New implementation of critical forwarding path as a “per-hop behavior”
 - Opportunities: leveraging Internet protocol base
 - Vendors: Opportunity for router upgrades
 - Small/medium-sized providers: economic necessity.
 - Large providers: view diff-serv as an intermediate solution to QoS while waiting for MPLS to integrate ATM, FR facilities and get traffic engineering features.

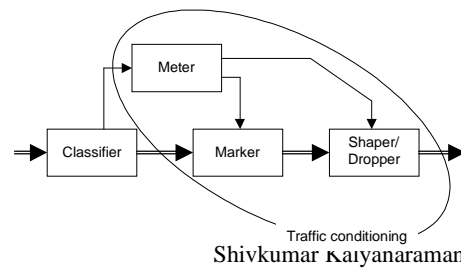
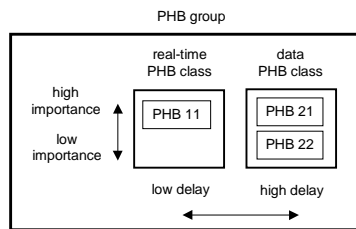
Differentiated Services Model



- **Network edge routers:** traffic conditioning (policing, marking, dropping), SLA negotiation
 - Set values in DS-byte based upon negotiated service and observed traffic. Per-flow state.
- **Interior routers:** traffic classification and forwarding
 - Use DS-byte as index into forwarding table

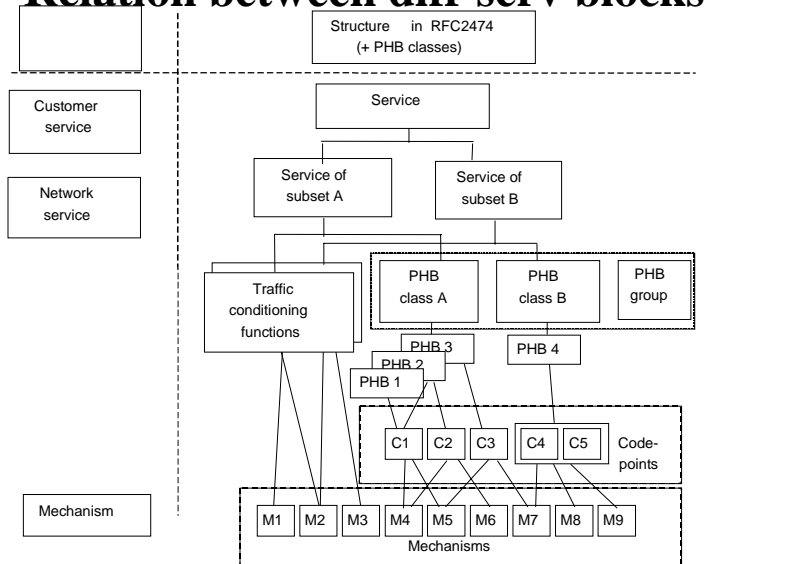
Diff-serv building blocks

- ❑ **Per-hop Behavior: (PHB)** generalization of mechanisms applied to a flow in the forwarding path
- ❑ **PHB Group:** Inter-related PHBs used together to implement a service.
- ❑ **Codepoints:** Bit combinations in the DS-byte
- ❑ **Mechanisms:** low level impln of building blocks
- ❑ **Traffic conditioners:** markers, meters, shapers etc



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Relation between diff-serv blocks



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IP Differentiated Services

- ❑ Only building blocks, no fully defined services
- ❑ Works with IPv4
- ❑ **Services:** leased-line emulation("premium service"), frame-relay emulation ("assured service"), CoS (Class-of-Service)
- ❑ Only data-plane building blocks defined: traffic conditioners, Per-hop Behaviors (PHBs)
- ❑ **Critique:** control-plane components undefined (contenders: RSVP, COPS, SNMP, MPLS, L2TP)

Control plane: MPLS

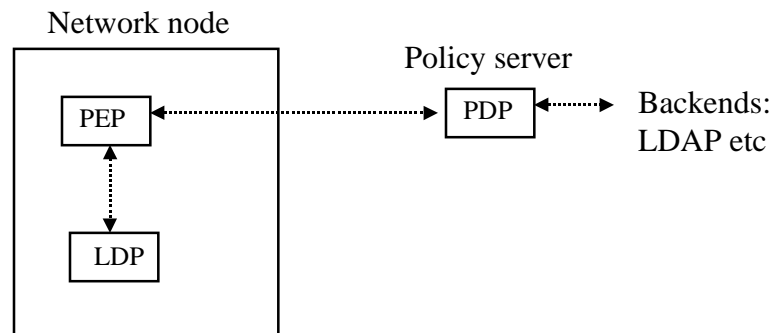
- ❑ Provides a framework for routing evolution
 - ❑ De-couples forwarding from routing control
 - ❑ Explicit routing
 - ❑ Constraint-based (QoS) routing, load-balancing
 - ❑ Traffic engineering: aggregating traffic flows into trunks, and mapping them onto pre-defined paths
- ❑ Provides a framework for integrating IP, ATM, and frame-relay cores
 - ❑ Allows re-engineering of the ATM control plane, and the IP forwarding plane

MPLS: building blocks

- ❑ ***Label***: short, fixed length field
- ❑ ***Forwarding table structure***:
 - ❑ Incoming label + subentry = outgoing label, outgoing interface, next-hop address (will include PHBs for diff-serv)
- ❑ ***Carrying label in header***:
 - ❑ Use VCI/VPI or DLCI in ATM or FR
 - ❑ New “shim” header for other link layers
- ❑ Forwarding algorithm: ***Label swapping***.
 - ❑ Use label as an index (exact match)
- ❑ ***Control component***:
 - ❑ Responsible for distributing routing & label-binding information: extensions to routing protocols, RSVP, LDP

COPS

- ❑ Common Open Policy Service
- ❑ Initially designed for adding ***policy control*** to RSVP
- ❑ Now being extended to support ***provisioning***
- ❑ Uses TCP; stateful exchange; common object model



Missing pieces in diff-serv

- ❑ **Provisioning/policy/signaling:** Assumed to be done using RSVP, COPS, SNMP, LDAP or over-engineering!
- ❑ **Route pinning/multi-paths:** extensions to OSPF, BGP, QoS routing
- ❑ **End-to-end services:** combination of above pieces: eg: frame-relay emulation, virtual leased line etc
- ❑ Tools to prevent traffic based denial of service attacks

Summary



- ❑ QoS big picture; ATM and IP building blocks/services
- ❑ Real-time transport/middleware: RTP, H.323
- ❑ Integrated services: RSVP, 2 services, scheduling, admission control etc
- ❑ Diff-serv: edge-routers, core routers; DS byte marking and PHBs
- ❑ Missing pieces: routing support (MPLS), pricing models, policy management (COPS)