

BOOTP, DHCP and NAT

Shivkumar Kalyanaraman
Rensselaer Polytechnic Institute
shivkuma@ecse.rpi.edu
<http://www.ecse.rpi.edu/Homepages/shivkuma>

Rensselaer Polytechnic Institute

Shivkumar Kalyanaraman

1



- Bootstrapping (Diskless workstations)
- BOOTP
- Dynamic address allocation
- DHCP
- Private Addresses: NAT, RSIP
- Ref: Chap 16, Doug Comer's TCP/IP book,
IETF NAT Working Group

Rensselaer Polytechnic Institute

Shivkumar Kalyanaraman

2

Bootstrapping

- Computer loads a simple boot program. The boot program loads operating system.
- On diskless machine, the computer needs to know the network address of the o/s file
- It needs to know its *own IP address., subnet mask, IP address of default router, IP address of DNS server*
- It only knows its h/w address.

Rensselaer Polytechnic Institute

Shivkumar Kalyanaraman

3

Configuration

- Different nodes have different parameters
- Configuration = Setting the parameters
- Key parameters for IP hosts:
 - IP Address
 - Default router address
 - Subnet mask
 - Name
 - DNS server IP addresses

Rensselaer Polytechnic Institute

Shivkumar Kalyanaraman

4

Key RARP Limitations

- RARP: user process over link layer *directly*
 - RARP server system-dependent
 - Needs to interface with link layer driver directly => separate filters and direct access to hardware needed
- Returns only IP address
 - Booting and configuration params not returned even though there is space in packet
 - Host *needs ICMP and TFTP to complete booting*
- Can't relay RARP requests to a central server.
 - Need RARP server per broadcast domain

Rensselaer Polytechnic Institute

Shivkumar Kalyanaraman

5

Method 2: BOOTP

- Runs over UDP/IP!
 - IP software can broadcast (to 255.255.255.255) even if local IP address unknown => client broadcasts BOOTP request
 - Port number 67 for server and 68 for client (not an ephemeral port)
 - Delivers BOOTP reply to BOOTP client and not other UDP apps when reply is broadcast
 - Does not wake up other servers during broadcast reply

Rensselaer Polytechnic Institute

Shivkumar Kalyanaraman

6

BOOTP (Continued)

- BOOTP requests/replies sent w/ DF bit set.
- Server can send reply via broadcast or unicast:
 - For unicast reply, BOOTP server knows the IP address, but the link layer address is not in the ARP cache
 - Note that the server cannot send an ARP message because client does not know its IP address
 - Server can use `ioctl(8)` (or `arp -s`) to set the value of the cache based upon BOOTP request => can do this only if it has permission

BOOTP Features (Continued)

- Else send broadcast reply
- Reply: IP Address, Boot Server IP address, Default Router, Boot file name, subnet mask
 - More information, but still only a single packet exchange
 - Client gets boot image using TFTP => booting still a 2-step process

BOOTP features (Continued)

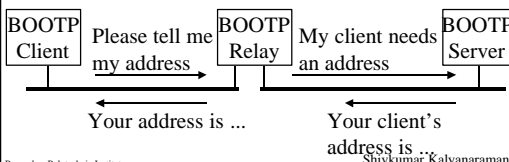
- Advantages of using UDP/IP:
 - Bootstrapping can occur across a router via a relaying mechanism
 - BOOTP uses checksum provided by UDP
- Multiple requests/replies
 - Process the first one
 - Client uses a transaction ID field to sort out replies
- Clients responsible for reliability
 - Uses timeout, retransmission & exponential backoff
 - Random initial timeout (betn 0 & 4s): simultaneous reboot after power restoration.

BOOTP Message Format

Operation	H/W Type	H/W Length	Hops	
Transaction Identifier				
Seconds elapsed		Unused		
Client IP Address				
Your IP Address				
Server IP Address				
Router IP Address				
Client H/W address				16 B
Server Host Name				64 B
Bootfile Name				128 B
Vendor Specific Area				64 B

BOOTP Message (Continued)

- Operation: 1 = Request, 2 = Reply
- H/w type: 1 = Ethernet
- H/w Address Length
- Hops: Initialized to zero. Incremented by BOOTP relays (routers)



BOOTP Message

- Boot File name: Generic name like "unix" in the request. Full name in response.
- Vendor specific area: Misnomer. Also used for general purpose info.
 - Magic cookie: First 4 octets = 99.130.83.99
 - Type-length-value: describes the option

Item	Code	Length
Padding	0	-
Subnet mask	1	4
Time of Day	2	4
End	255	-

DHCP

- BOOTP limitation: cannot dynamically assign IP address
- Dynamic Host Configuration Protocol (DHCP)
 - BOOTP + Dynamic allocation of IP addresses => compatible with BOOTP.
 - No new fields in header.
 - Addresses can be leased for a period. Reallocated to the same or other nodes after lease expiry.

DHCP Message Format

Operation	H/W Type	H/W Length	Hops
Transaction Identifier			
Seconds elapsed		Flags	
Client IP Address			
Your IP Address			
Server IP Address			
Router IP Address			
Client H/W address			
Server Host Name			
Bootfile Name			
Options (Variable)			

DHCP Message Format

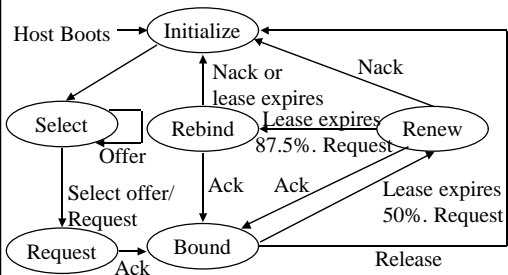
- Slightly modified version of BOOTP message => **A DHCP server can be programmed to answer BOOTP requests**
- BOOTP's Unused field renamed to :Flags"
- Only one bit of 16-bit Flags has been defined
 - **Left-most flag bit =1 => Servers, please reply using IP broadcast address**
- Servers by default send hardware unicast response
- Vendor-specific field renamed to "Options"
 - **Size increased to 312 bytes (from 64 bytes)**
 - **Option type 53 specifies the "type of the message"**

DHCP (Continued)

Type	Meaning
1	DHCP Discover
2	DHCP Offer
3	DHCP Request
4	DHCP Decline
5	DHCP Ack
6	DHCP Nack
7	DHCP Release

- "Option overload"
 - Server Host name and boot file name when unused for their original purpose could be used to code more options

DHCP State Diagram



DHCP States

- Boots => INITIALIZE state
- DHCPDISCOVER: broadcast request to servers => SELECT state
- DHCPOFFER (from server) => remain in SELECT
- DHCPREQUEST => select one of the offers and notify server (goto REQUEST state) about the lease

DHCP States (Continued)

- DHCPACK => server OKs request to lease => go to the BOUND state
- Renewal: after 50% of lease go to RENEW state
- Rebind: after 87.5% of time, if server has not responded, try again and go to REBIND.
- If server NACKs or lease expires, or client sends DHCPRELEASE, go to INITIALIZE, else come back to BOUND state

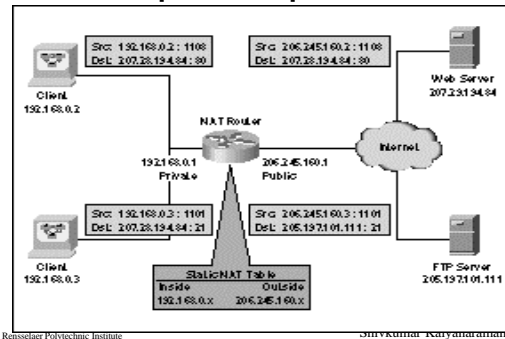
Private vs Public Addresses

- Since IPv4 addresses are scarce, enterprises may use private addresses within their "realms"

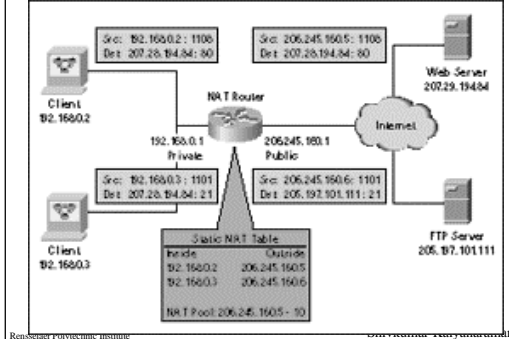
Class	Private Address Range
A	10.0.0.0 ... 10.255.255.255
B	172.16.0.0 ... 172.16.255.255
C	192.168.0.0 ... 192.168.255.255

- Need to get "globally unique" public addresses for external use.
- Mapping between public & private addresses done by NAT (Network Address Translator)

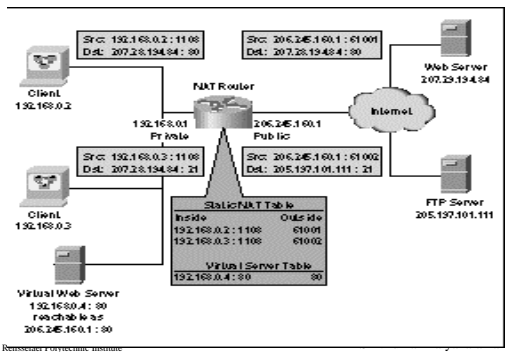
Simple NAT operation



Dynamic NAT = NAT + DHCP



Network Address Port Translation (NAPT)



NAPT (contd)

- Also known as IP masquerading. Allows many hosts to share a single IP address differentiated by port numbers.
- Eg: Suppose private hosts 192.168.0.2 and 192.168.0.3 send packets from source port 1108.
- NAPT translates these to a single public IP address 206.245.160.1 and two different source ports, say 61001 and 61002.
 - Response traffic received for port 61001 is routed back to 192.168.0.2:1108,
 - Traffic for port 61002 traffic is routed back to 192.168.0.3:1108.

Realm-Specific IP (RSIP)

- NAT (and NAPT) have to mess with several transport/application level fields.
- NAT breaks IPSec.... Solution: RSIP
- RSIP leases public IP addresses and ports to RSIP hosts => not transparent like NAT.
 - RSIP does not operate in stealth mode and does not translate addresses on the fly.
 - RSIP allows hosts to directly participate concurrently in several addressing realms.
 - Avoids violating the end- to-end nature of the Internet => allows IPSec

Summary



- RARP allows finding an IP address
- BOOTP allows default router, subnet mask, DNS
- DHCP allows dynamic allocation
- DHCP is backward compatible with BOOTP
- NAT, NAPT, RSIP allow use of private addresses and smaller pool of public addresses