

IP Next Generation (IPv6)

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- Limitations of current Internet Protocol (IP)
- How many addresses do we need?
- IPv6 Addressing
- IPv6 header format
- IPv6 features: routing flexibility, plug-n-play, multicast support, flows

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IP Addresses

- **Example:** 164.107.134.5
= 1010 0100 : 0110 1011 : 1000 0110 : 0000 0101
= A4:6B:86:05 (32 bits)
- Maximum number of address = 2^{32} = 4 Billion
- Class A Networks: 15 Million nodes
- Class B Networks: 64,000 nodes or less
- Class C Networks: 250 nodes or less

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IP Address Format

- Three all-zero network numbers are reserved
- 127 Class A + 16,381 Class B + 2,097,151 Class C networks = 2,113,659 networks total
- Class B is most popular.
- 20% of Class B were assigned by 7/90 and doubling every 14 months \Rightarrow Will exhaust by 3/94
- Question: Estimate how big will you become?
Answer: More than 256!
Class C is too small. Class B is just right.

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How Many Addresses?

- 10 Billion people by 2020
- Each person has more than one computer
- Assuming 100 computers per person $\Rightarrow 10^{12}$ computers
- More addresses may be required since
 - Multiple interfaces per node
 - Multiple addresses per interface
 - Some believe 2^6 to 2^8 addresses per host
- Safety margin $\Rightarrow 10^{15}$ addresses
- IPng Requirements $\geq 10^{12}$ end systems and 10^9 networks. Desirable 10^{12} to 10^{15} networks

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How big an address space ?

- H Ratio = $\log_{10}(\# \text{ of objects})/\text{available bits}$
- 2^n objects with n bits: H-Ratio = $\log_{10}2 = 0.30103$
- French telephone moved from 8 to 9 digits at 10^7 households $\Rightarrow H = 0.26$ (~3.3 bits/digit)
- US telephone expanded area codes with 10^8 subscribers $\Rightarrow H = 0.24$
- Physics/space science net stopped at 15000 nodes using 16-bit addresses $\Rightarrow H = 0.26$
- 3 Million Internet hosts currently using 32-bit addresses $\Rightarrow H = 0.20 \Rightarrow$ A few more years to go

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IPv6 Addresses

- 128-bit long. Fixed size
- $2^{128} = 3.4 \times 10^{38}$ addresses
 - ⇒ 665×10^{21} addresses per sq. m of earth surface
- If assigned at the rate of $10^6/\mu\text{s}$, it would take 20 years
- Expected to support 8×10^{17} to 2×10^{33} addresses
 - $8 \times 10^{17} \Rightarrow 1,564$ address per sq. m

IPv6 Addresses (Continued)

- Allows multiple interfaces per host.
- Allows multiple addresses per interface
- Allows unicast, multicast, anycast
- Allows provider based, site-local, link-local
- 85% of the space is unassigned

Colon-Hex Notation

- **Dot-Decimal:** 127.23.45.88
- **Colon-Hex:**
FEDC:0000:0000:0000:3243:0000:0000:ABCD
 - Can skip leading zeros of each word
 - Can skip one sequence of zero words, e.g.,
FEDC::3243:0000:0000:ABCD or
::3243:0000:0000:ABCD
 - Can leave the last 32 bits in dot-decimal, e.g.,
::127.23.45.88
 - Can specify a prefix by /length, e.g.,
2345:BA23:7::/40

Header

- **IPv6:**

Version	Class	Flow Label	
Payload Length		Next Header	Hop Limit
Source Address			
Destination Address			

- **IPv4:**

Version	IHL	Type of Service	Total Length	
Identification		Flags	Fragment Offset	
Time to Live	Protocol	Header Checksum		
Source Address				
Destination Address				
Options				Padding

IPv6 vs IPv4

- IPv6 only twice the size of IPv4 header
- Only version number has the same position and meaning as in IPv4
- **Removed:**
 - Header length, fragmentation fields (identification, flags, fragment offset), header checksum
- **Replaced:**
 - Datagram length by payload length
 - Protocol type by next header
 - Time to live by hop limit
 - Type of service by "class" octet
- **Added:** flow label
- All fixed size fields.

IPv6 vs IPv4 (Continued)

- No optional fields. Replaced by extension headers.
 - Idea: avoid unnecessary processing by intermediate routers while not sacrificing the flexibility possible due to options
 - Next Header = 6 (TCP), 17 (UDP), etc

Extension Headers

Base Header	Extension Header 1	Extension Header n	Data
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- Most extension headers are examined only at destination
- Routing: Loose or tight source routing
- Fragmentation: All IPv6 routers can carry 536 Byte payload
- Authentication
- Hop-by-Hop Options
- Destination Options:

Extension Header (Continued)

- Only Base Header:

Base Header	TCP Segment
Next = TCP	

- Only Base Header and One Extension Header:

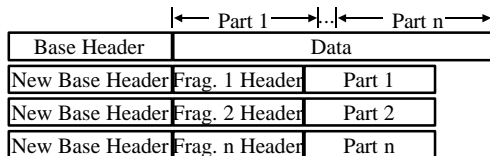
Base Header	Route Header	TCP Segment
Next = TCP	Next = TCP	

- Only Base Header and Two Extension Headers:

Base Header	Route Header	Auth Header	TCP Segment
Next = TCP	Next = Auth	Next = TCP	

Fragmentation

- Routers cannot fragment. Only source hosts can.
 - ⇒ Need path MTU discovery or tunneling
- Fragmentation requires an extension header
- Payload is divided into pieces
- A new base header is created for each fragment



IPv6 Addressing and Routing

- Aggregatable Global Unicast Addresses
- Link-local and Site-local addresses
- Multicast and Anycast support
- Provider-based inter-domain routing & IDRP

Initial IPv6 Prefix Allocation

Allocation	Prefix	Allocation	Prefix
Reserved	0000 0000	Unassigned	101
Unassigned	0000 0001	Unassigned	110
NSAP	0000 001	Unassigned	1110
IPX	0000 010	Unassigned	1111 0
Unassigned	0000 011	Unassigned	1111 10
Unassigned	0000 1	Unassigned	1111 110
Unassigned	0001	Unassigned	1111 1110
Unassigned	001	Unassigned	1111 1110 0
Provider-based*	010	Link-Local	1111 1110 10
Unassigned	011	Site-Local	1111 1110 11
Geographic	100	Multicast	1111 1111

*Has been renamed as "Aggregatable global unicast"

Aggregatable Global Unicast Addresses

- Address allocation: "provider-based" plan
- Format: TLA + NLA + SLA + 64-bit interface ID
- TLA = "Top level aggregator." (13 bits)
 - Ranges of TLA values allocated to various registries
 - For "backbone" providers or "exchange points"

Aggregatable Global Unicast Addresses (Continued)

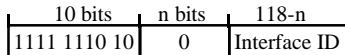
- NLA = "Next Level Aggregator" (32 bits)
 - Second tier provider and a subscriber
 - More levels of hierarchy possible within NLA
- SLA = "Site level aggregator" = 16 bits for link
 - Renumbering after change of provider => change the TLA and NLA. But have same SLA & I/f ID

Aggregatable Global Unicast Addresses (Continued)

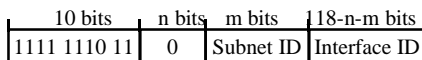
- Interface ID = 64 bits
 - Will be based on IEEE EUI-64 format
 - An extension of the IEEE 802 (48 bit) format.
 - Possible to derive the IEEE EUI-64 equivalent of current IEEE 802 addresses
 - Along with neighbor discovery procedures, obviates need for ARP.

Local-Use Addresses

- **Link Local:** Not forwarded outside the link, FE:80::xxx

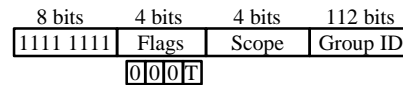


- **Site Local:** Not forwarded outside the site, FE:C0::xxx



- Provides plug and play

Multicast Addresses



- All routers recognize this format
 - all routers can route multicast packets.
 - Also IGMP part of ICMPv6 => required.
 - I.e. IPv6 supports native multicast
- **Flags:** T = 0 => Permanent (well-known) multicast address, 1 => Transient

Multicast Addresses (Continued)

- **Scope:** 1 Node-local, 2 Link-local, 5 Site-local, 8 Organization-local, E Global => routers reqd to honor this.
- **Predefined:** 1 => All nodes, 2 => Routers, 1:0 => DHCP servers

Multicast & Anycast

- **Scoping.** Eg: 43 => NTP Servers
 - FF01::43 => All NTP servers on this node
 - FF02::43 => All NTP servers on this link
 - FF05::43 => All NTP servers in this site
 - FF08::43 => All NTP servers in this org.
 - FFOF::43 => All NTP servers in the Internet
- **Structure of Group ID:**
 - First 80 bits = zero (to avoid risk of group collision, because IP multicast mapping uses only 32 bits)

Address Auto-configuration

- Allows plug and play
- BOOTP and DHCP are used in IPv4
- DHCPng will be used with IPv6
- Two Methods: Stateless and Stateful
- Stateless:
 - A system uses *link-local* address as source and multicasts to "All routers on this link"
 - Router replies and provides all the needed prefix info

Address Auto-configuration (Continued)

- All prefixes have a associated *lifetime*
- System can use link-local address permanently if no router
- Stateful:
 - Problem w stateless: Anyone can connect
 - Routers ask the new system to go DHCP server (by setting managed configuration bit)
 - System multicasts to "All DHCP servers"
 - DHCP server assigns an address

ICMPv6: Neighbor Discovery

- ICMPv6 combines regular ICMP, ARP, Router discovery and IGMP.
- The "neighbor discovery" is a generalization of ARP & router discovery.
- Source maintains several caches:
 - destination cache: dest -> neighbor mapping
 - neighbor cache: neighbor IPv6 -> link address
 - prefix cache: prefixes learnt from router advertisements
 - router cache: router IPv6 addresses

Neighbor Discovery (Continued)

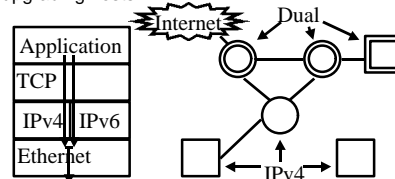
- Old destination => look up destination cache
- If new destination, match the prefix cache. If match => destination local!
- Else select a router from router cache, use it as the next-hop (neighbor).
 - Add this neighbor address to the destination cache
- Solicitation-advertisement model:
 - Multicast solicitation for neighbor media address if unavailable in neighbor cache
 - Neighbor advertisement message sent to soliciting station.

Real-Time/QoS Support

- Flow label and the "class" octet field
- Flow = sequence of packets from a single source to a particular (unicast/multicast) destinations requiring special handling by intermediate routers
 - Class field can simplify packet classification (data-plane support for QoS)
 - "Class" field assignments being worked upon by differentiated services group at IETF

Transition Mechanisms

- Dual-IP Hosts, Routers, Name servers
- Tunneling IPv6 over IPv4
- Nodes can be partially upgraded to IPv6
- It is better (though not required) to upgrade routers before upgrading hosts



Application Issues

- Most application protocols will have to be upgraded: FTP, SMTP, Telnet, Rlogin
- 27 of 51 Full Internet standards, 6 of 20 draft standards, 25 of 130 proposed standards will be revised for IPv6
- No checksum ⇒ checksum at upper layer is mandatory, even in UDP
- non-IETF standards: X-Open, Kerberos, ... will be updated
- Should be able to request and receive new DNS records

Summary



- IPv6 uses 128-bit addresses
- Allows provider-based, site-local, link-local, multicast, anycast addresses
- Fixed header size. Extension headers instead of options for provider selection, security etc
- Allows auto-configuration
- Dual-IP router and host implementations for transition