


TCP (Part III: Misc)

Shivkumar Kalyanaraman
Rensselaer Polytechnic Institute
shivkuma@ecse.rpi.edu
<http://www.ecse.rpi.edu/Homepages/shivkuma>

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Overview

- ❑ TCP Persist and Keepalive timers
- ❑ Silly window syndrome
- ❑ Path MTU
- ❑ Window Scale Factor
- ❑ Timestamp option
- ❑ T/TCP: TCP for transactions

Ref: Chap 22, 23, 24; RFC 1323

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TCP Persist Timer

- ❑ Receiver flow control can set window to zero
- ❑ Receiver later sends "window update acks"
- ❑ But TCP does not transmit acks reliably => update acks may be lost and source may be stuck at a zero window value
- ❑ TCP uses persist timer to query the receiver periodically to find if the window has been increased.
- ❑ Persist timer always bounded between 5s and 60s. It does exponential backoff like other timers too.

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Silly Window Syndrome

- ❑ A) The system operates at a small window (sends segments which are not MSS-sized) even if the receiver grants a large window.
- ❑ B) Receiver advertises small windows.
- ❑ Solution: **batching**
 - ❑ Receiver must not advertise small windows
 - ❑ **Sender waits until segment full before sending** (extension of Nagle's algo),
 - ❑ It can transmit everything if it is not waiting for any ACK (or if Nagle's algo has been disabled)

TCP Keepalive timer

- ❑ Optional timer.
- ❑ Not part of TCP spec, but found in most implementations.
 - ❑ Not necessary, because "connection" defined by endpoints.
 - ❑ Connection can be "up" as long as source/destination "up".
- ❑ Typical use: to detect idle clients or half-open connections and de-allocate server resources tied up to them. Eg: telnet, ftp.

Path MTU discovery

- ❑ Assume $MSS = Min(\text{local MTU} - \text{headers}, \text{destination MSS})$. Set DF bit.
- ❑ If ICMP error, reduce segment size and retransmit.
- ❑ Since routes change dynamically, a larger value can be tried again after a time interval (RFC 1191 recommends 10 min, but Solaris uses 30 s).

Gigabit Networks

- ❑ “Higher **Bandwidth** Networks”
- ❑ Propagation latency unchanged.
 - ❑ Increasing bandwidth from 1.5Mb/s to 45 Mb/s (factor of 29) decreases file transfer time of 1MB by a factor of 25.
 - ❑ But, increasing from 1 Gb/s to 2 Gb/s gives an **improvement of only 10% !**
 - ❑ Transfer time = propagation time + transmission time + queueing/processing.
- ❑ Design networks to minimize delay (queueing, processing, reduce retransmission latency)


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Window Scaling Option

- ❑ Long Fat Pipe Networks (LFN): Satellite links
- ❑ Need very large window sizes.
- ❑ Normally, Max window = $2^{16} = 64$ KBytes
- ❑ **Window scale:** Window = $W \times 2^{\text{Scale}}$

Kind = 3	Length = 3	Scale
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- ❑ Max window = $2^{16} \times 2^{255}$
- ❑ Option sent only in SYN and SYN + Ack segments.
- ❑ RFC 1323



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Timestamp option

- ❑ For LFNs, need accurate and more frequent RTT estimates.
- ❑ Timestamp option:
 - ❑ Place a timestamp value in any segment.
 - ❑ Receiver echoes timestamp value in ack
 - ❑ If acks are delayed, the timestamp value returned corresponds to the *earliest* segment being acked.
- ❑ Segments lost/retransmitted => RTT overestimated

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PAWS: Protection against wrapped sequence numbers

- ❑ Largest receiver window = $2^{30} = 1 \text{ GB}$
- ❑ “Lost” segment may reappear before MSL, and the sequence numbers may have wrapped around
- ❑ The receiver considers the timestamp as an extension of the sequence number => discard out-of-sequence segment based on both seq # and timestamp.
- ❑ Req: timestamp values need to be monotonically increasing, and need to increase by at least one per window


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T/TCP: Transaction Oriented TCP

- ❑ Three-way handshake => Long delays for transaction-oriented applications.
 - ❑ T/TCP *extension* avoids 3-way handshakes
 - ❑ Request/reply data sent with connection messages
 - ❑ Server caches a connection count (CC) per-client to detect duplicate requests and avoid replaying transaction
 - ❑ TIME_WAIT is shortened by setting it to $8 \cdot \text{RTO}$
 - ❑ Latency = RTT + server processing time (SPT)

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Summary



- ❑ Persist and keepalive timers, silly window avoidance
- ❑ Enhancements for LFNs: window scale option, timestamp option, PAWS
- ❑ T/TCP extension to TCP for transactions

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