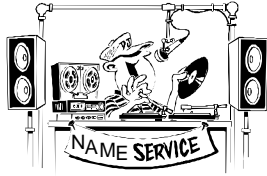


# Domain Name System (DNS)



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- Naming hierarchy
- Server hierarchy
- Name resolution
- Other information in name servers
- Ref: Chap 14

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## Why Names?

- Computers use addresses
- Humans cannot remember IP addresses  
⇒ Need names  
Example, "shiva" for 128.113.50.56
- Simplest Solution: Each computer has a unique name and has a built in table of name to address translation
- Problem: Not scalable
- Solution: DNS (Adopted in 1983)
- Hierarchical Names: shiva.ecse.rpi.edu

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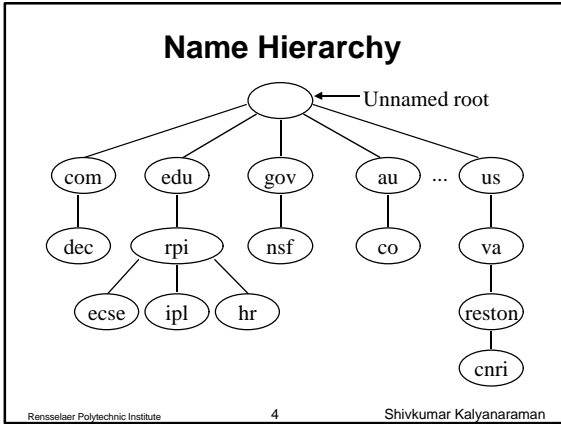
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- ### Name Hierarchy
- Unique domain suffix is assigned by Internet Authority
  - The domain administrator has complete control over the domain
  - No limit on number of subdomains or number of levels
  - computer.site.division.company.com
  - computer.site.subdivision.division.company.com
  - Domains within an organization do not have to be uniform in number of subdomains or levels
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- ### Name Hierarchy (Cont)
- Name space is not related to physical interconnection, e.g., ecse.rpi.edu and ipl.rpi.edu could be on the same floor or in different cities
  - Geographical hierarchy is also allowed, e.g., cnri.reston.va.us
  - A name could be a *subdomain* (eg: ecse.rpi.edu) or an individual *object* (eg: cortez.rpi.edu)
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## Top Level Domains

Domain Name/Assignment	
com	Commercial
edu	Educational
gov	Government
mil	Military
net	Network
org	Other organizations
arpa	Advanced Research Project Agency
country code	au, uk, ca

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## Server Hierarchy

- Servers are organized in a hierarchy
- Each server has an authority over a part of the naming hierarchy
- The server does not need to keep all names.
- It needs to know other servers who are responsible for other subdomains
- A single node in the naming tree cannot be split among multiple servers
- A given level of hierarchy can be partitioned into multiple servers

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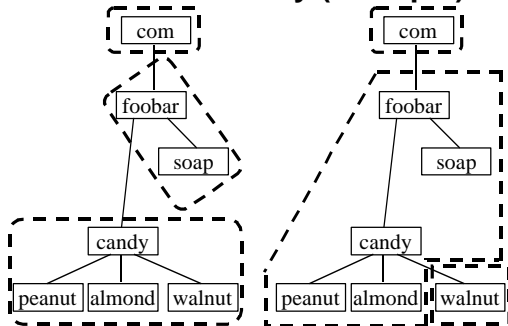
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## Server Hierarchy (example)



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## Server Hierarchy (Cont)

- ❑ Authority ⇒ has the name to address translation table
- ❑ Responsible ⇒ Either has the name to address translation table or knows the server who has
  - ❑ But such a reply is called “non-authoritative” reply
- ❑ A single server can serve multiple domains, e.g., purdue.edu and laf.in.us
- ❑ Root server knows about servers for top-level domains, e.g., com
- ❑ Each server knows the root server

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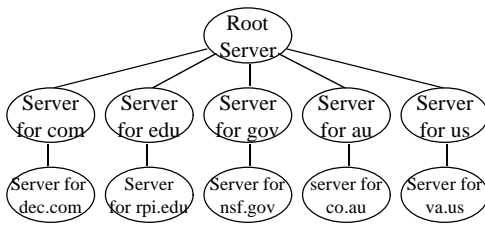
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## Server Hierarchy: Example



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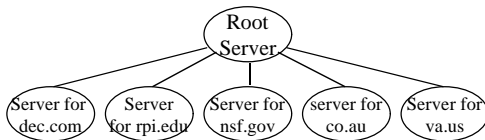
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## Server Hierarchy: Better



- ❑ Fewer servers
- ❑ More entries/links per server
- ❑ Fewer levels to traverse before resolving a name

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## Name Resolution (Cont)

- ❑ Each computer has a name resolver routine, e.g., `gethostbyname` & `gethostbyaddr` in UNIX
- ❑ Each resolver knows the name of a local DNS server
- ❑ Resolver sends a DNS request to the server
- ❑ DNS server either gives the answer, forwards the request to another server, or gives a referral
- ❑ Referral = Next server to whom request should be sent

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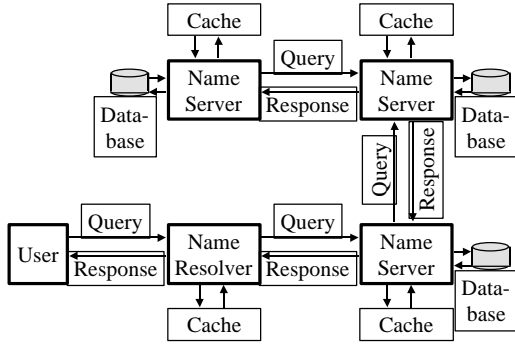
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## Name Resolution



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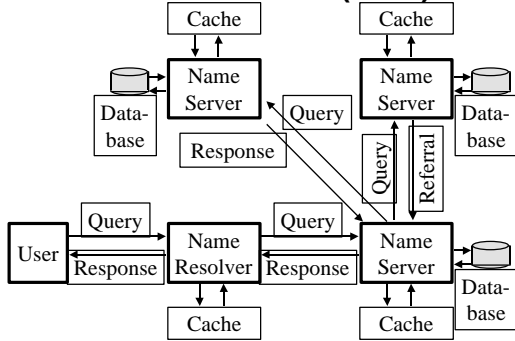
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## Name Resolution (Cont)



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## Name Resolution (Cont)

- Resolvers use UDP (single name) or TCP (whole group of names)
- Knowing the address of the root server is sufficient
- Recursive Query:  
Give me an answer (Don't give me a referral)
- Iterative Query:  
Give me an answer or a referral to the next server
- Resolvers use recursive query.
- Servers use iterative query.

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## DNS Optimization

- Spatial Locality: Local computers referenced more often than remote
- Temporal Locality: Same set of domains referenced repeatedly  $\Rightarrow$  Caching
- Each entry has a time to live (TTL)
- Replication: Multiple servers. Multiple roots. Ask the geographically closest server.

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## Abbreviations

- Servers respond to a full name only
- However, humans may specify only a partial name
- Resolvers may fill in the rest of the suffix, e.g., shiv = shiv.ecse.ohio-state.edu
- Each resolver has a list of suffixes to try

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## DNS Message Format

Identification	Parameter
Number of Questions	Number of Answers
Number of Authority	Number of Additional
Question Section	
...	
Answer Section	
...	
Authority Section	
...	
Additional Information Section	
...	

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## Format (Cont)

- **Format of the query section entries:**

Query Domain Name	
...	
Query Type	Query Class

- **Format of other section entries:**

Resource Domain Name	
Type	Class
Time to live	
Resource Data Length	Resource Data

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## DNS Message Format

- **Length = 0 ⇒ End of names. Length < 64  
Two msbs (most significant bits) = 11 ⇒  
Pointer**
- **Resource data contains serial (version)  
number of the zone, refresh interval, retry  
interval, expiry interval, mailbox of the  
responsible person, etc.**

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## DNS Message Format (Cont)

Bit	Meaning
0	Operation: 0=Query, 1=Response
1-4	Query type: 0=Standard, 1=Inverse, 2,3 obsolete
5	Set if answer authoritative
6	Set if message truncated
7	Set if recursion desired
8	Set if recursion available
9-11	Reserved
12-15	Response type: 0=No error, 1=Format error, 2=Server Failure, 3=Name does not exist

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## Types of DNS Entries

- DNS used other types of resolution
- Eg: also for finding mail server, pop server, responsible person, etc for a computer
- DNS database has multiple "types"
  - Record type A ⇒ Address of X
  - Record type MX ⇒ Mail exchanger of X
- DNS database may also have multiple "classes"
  - Can support name resolution for multiple protocols eg: IP, SNA, DECbit etc
- Pointer queries: given IP address find name

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## Resource Record Types

Type	Meaning
A	Host Address
CNAME	Canonical Name (alias)
HINFO	CPU and O/S
MINFO	Mailbox Info
MX	Mail Exchanger
NS	Authoritative name server for a domain
PTR	Pointer to a domain name (link)
RP	Responsible person
SOA	Start of zone authority (Which part of naming hierarchy implemented)
TXT	Arbitrary Text

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## Summary



- ❑ **DNS: Maps names to addresses**
- ❑ **Names are hierarchical. Administration is also hierarchical.**
- ❑ **No standard for number of levels**
- ❑ **Replication and caching is used for performance optimization.**

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## Informal Exercises

- ❑ **/etc/hosts is a table of name-to-IP-address mappings**
  - ❑ **Find out approximately how many hosts, subnets and domains are there in the RPI campus net**
  - ❑ **Does this table give some addresses of root servers as well ?**
  - ❑ **Why does the /etc/hosts in ECSE have the entire table for the campus net.**
- ❑ **Lookup the file /etc/resolv.conf which gives the domain name and addresses of nameservers.**
  - ❑ **Why are multiple nameservers listed ?**
  - ❑ **Lookup the name of an IP address using nslookup. This generates a pointer query - you can watch it using tcpdump. What is unusual about pointer queries on the wire ?**

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