Internet Protocols: Quiz 1

- □ This quiz consists of true/false questions for 25 pts and two quantitative problems for 25 pts.
- □ In the True/False questions, the following grading policy will be used:
 - □ Correct answer: +1 pt
 - □ Wrong answer: -1 pt (*negative grading is used*)
 - **Blank/Unattempted: 0 pts**
- □ There will be *no negative grading for the quantitative problems*. Partial credit may be awarded where appropriate.
- Open book policy
- □ *Time: 45 min.* Strictly enforced.
- □ This is the first quiz out of three quizzes. *Best two out of three* will be considered for final grades. Each of the two quizzes chosen will be weighted equally.

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True or False? (25 points)

Note: Correct Ans = +1; Wrong Answer = -1; Did not attempt = 0

ΤF

- Layering is desired because it is the most efficient way of designing and implementing network protocols
- □ □ If a router looks at the TCP or UDP port numbers to base any of its decisions, it is a violation of layering.
- □ □ A multihomed host must be configured as a router to allow communication between the networks on the two interfaces
- □ □ The sockets API models the network as an I/O device with the open-read-writeclose paradigm, the difference being that a socket need not be "bound" to an address upon creation.
- □ □ Typically, ISPs assign IP addresses dynamically to its dial-up clients
- □ □ As a packet passes from one end to another, it will change some of its address fields depending upon the network it traverses
- □ □ Ethernet and IP perform a limited protocol-based demultiplexing, whereas TCP/UDP ports allow more flexible port-based demultiplexing

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ΤF

- □ □ A telnet server demultiplexes incoming TCP segments based upon its local IP address and port number.
- \Box \Box A collision domain marks the boundaries of an Ethernet LAN
- □ □ The 48-bit LAN address has internal structure, but it is considered a "flat" address since the entire address is required at every stage to forward the packet
- □ □ The key difference between Ethernet and 802.3 is that the latter has a length field, which means that the former cannot support variable length packets.
- □ □ Typical IP overhead is 20 bytes while Ethernet overhead is 14 bytes
- □ □ The Initial Seq Number (ISN) is periodically incremented to avoid confusion from packets belonging to previous incarnations
- □ □ SLIP and PPP both support dynamic IP address assignment
- □ □ When a header checksum error is detected, IP quietly drops the packet and reports the error to the source
- □ □ Fragments are created at 8-octet boundaries

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True or False?

ΤF

- □ □ A result of the "end-to-end" principle was that complex control functions were pushed to the edge while the forwarding path was kept as simple as possible.
- □ □ Subnetting allows more levels of hierarchy in the addressing structure.
- □ □ The IP addresses 128.40.30.20 and 128.40.30.45 belong to the same subnet
- □ □ Subnetting transforms classful addressing into classless addressing
- □ □ The reason IP addressing is hierarchical is because the router can look at a portion of the address to decide where to forward it.
- □ □ Though the IP max length is 65535 octets, a destination need not accept a datagram larger than 576 bytes
- □ □ If a UDP checksum value is zero, it means that the sender did not compute a checksum
- \Box \Box On an Ethernet, the MSS is 1500 bytes
- □ □ The 2MSL wait allows TCP servers to be brought down and brought up immediately

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1) a) (7 pts) The IP checksum involves 1's complement arithmetic on 16-bit quantities. Use a similar technique, but on 4-bit quantities to compute the blank checksum field:

1111 0000 1100 ____ 0101 1000

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2) a) (13 pts) An IP datagram of length 2000 bytes needs to cross an Ethernet (MTU = 1500B) followed by a WAN (MTU = 576B). How many fragments reach the destination ? What are the values of the Header length, More bit, Offset, and Length fields in each fragment ?

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True or False? (25 points)

Note: Correct Ans = +1; *Wrong Answer* = -1; *Did not attempt* = 0

ΤF

- $\Box \sqrt{}$ Layering is desired because it is the most efficient way of designing and implementing network protocols
- $\sqrt{\Box}$ If a router looks at the TCP or UDP port numbers to base any of its decisions, it is a violation of layering.
- $\Box \sqrt{A}$ Multihomed host must be configured as a router to allow communication between the networks on the two interfaces
- $\sqrt{\Box}$ The sockets API models the network as an I/O device with the open-read-writeclose paradigm, the difference being that a socket need not be "bound" to an address upon creation.
- $\sqrt{\Box}$ Typically, ISPs assign IP addresses dynamically to its dial-up clients
- $\sqrt{\Box}$ As a packet passes from one end to another, it will change some of its address fields depending upon the network it traverses
- $\sqrt{\Box}$ Ethernet and IP perform a limited protocol-based demultiplexing, whereas TCP/UDP ports allow more flexible port-based demultiplexing

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ΤF

- $\Box \sqrt{A}$ telnet server demultiplexes incoming TCP segments based upon its local IP address and port number.
- $\Box \sqrt{A}$ A collision domain marks the boundaries of an Ethernet LAN
- $\sqrt{\Box}$ The 48-bit LAN address has internal structure, but it is considered a "flat" address since the entire address is required at every stage to forward the packet
- $\Box \sqrt{}$ The key difference between Ethernet and 802.3 is that the latter has a length field, which means that the former cannot support variable length packets.
- $\Box \sqrt{10}$ Typical IP overhead is 20 bytes while Ethernet overhead is 14 bytes
- $\sqrt{\Box}$ The Initial Seq Number (ISN) is periodically incremented to avoid confusion from packets belonging to previous incarnations
- $\Box \sqrt{\text{SLIP}}$ and PPP both support dynamic IP address assignment
- $\Box \sqrt{}$ When a header checksum error is detected, IP quietly drops the packet and reports the error to the source
- $\sqrt{\Box}$ Fragments are created at 8-octet boundaries

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True or False?

ΤF

- $\sqrt{\Box}$ A result of the "end-to-end" principle was that complex control functions were pushed to the edge while the forwarding path was kept as simple as possible.
- $\sqrt{\Box}$ Subnetting allows more levels of hierarchy in the addressing structure.
- $\Box \sqrt{}$ The IP addresses 128.40.30.20 and 128.40.30.45 belong to the same subnet
- $\Box \sqrt{}$ Subnetting transforms classful addressing into classless addressing
- $\sqrt{\Box}$ The reason IP addressing is hierarchical is because the router can look at a portion of the address to decide where to forward it.
- $\sqrt{\Box}$ Though the IP max length is 65535 octets, a destination need not accept a datagram larger than 576 bytes
- $\sqrt{\Box}$ If a UDP checksum value is zero, it means that the sender did not compute a checksum
- $\Box \sqrt{}$ On an Ethernet, the MSS is 1500 bytes
- $\Box \sqrt{}$ The 2MSL wait allows TCP servers to be brought down and brought up immediately

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1) a) (7 pts) The IP checksum involves 1's complement arithmetic on 16-bit
quantities. Use a similar technique, but on 4-bit quantities to compute the
   blank checksum field:
       1111 0000 1100 ____ 0101 1000
Checksum = 1s complement sum of the 1s complement of 4-bit quantities.
1s complement of 1111, 0000, 1100, 0101, 1000
             is 0000, 1111, 0011, 1010, 0111.
1s complement sum: 0000 + 1111 = 1111.
                    1111 + 0011 = 0010 + 1 (carry) = 0011
                    0011 + 1010 = 1101
                    1101 + 0111 = 0100 + 1 (carry) = 0101
Ans: Checksum = 0101
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2) a) (13 pts) An IP datagram of length 2000 bytes needs to cross an Ethernet (MTU = 1500B) followed by a WAN (MTU = 576B). How many fragments reach the destination ? What are the values of the More bit, (fragment) offset, and Length fields in each fragment ?

IP Datagram 2000B => payload = 1980B > Enet MTU = 1500B

- => Max IP payload is nearest multiple of 8 to 1480B (1500B 20B) = 1480B
- \Rightarrow 1st fragment: Length = (1480B + 20B) = 1500B; MF set; Fragoff = 0

2*nd* fragment: Length = (500B + 20B) = 520B; MF not set;

Fragoff (13-bit quantity) = 1480 >> 3 = 185

- WAN MTU = $576B \Rightarrow 1st$ fragment needs to be refragmented. Nearest multiple of 8 to (576B 20B = 556B) is 552B.
 - => Fragment 1a: Length = (552B + 20B) = 572B; MF set; Fragoff = 0 Fragment 1b: Length = (552B + 20B) = 572B; MF set; Fragoff = 69 Fragment 1c: Length = (376B + 20B) = 396B; MF set; Fragoff = 138 Fragment 2 not fragmented further.

Ans: Four fragments reach the destination with the fields highlighted above.

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