

Reference: IS-IS vs OSPF

Shivkumar Kalyanaraman
Rensselaer Polytechnic Institute
shivkuma@ecse.rpi.edu

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Shivkumar Kalyanaraman

IS-IS Overview

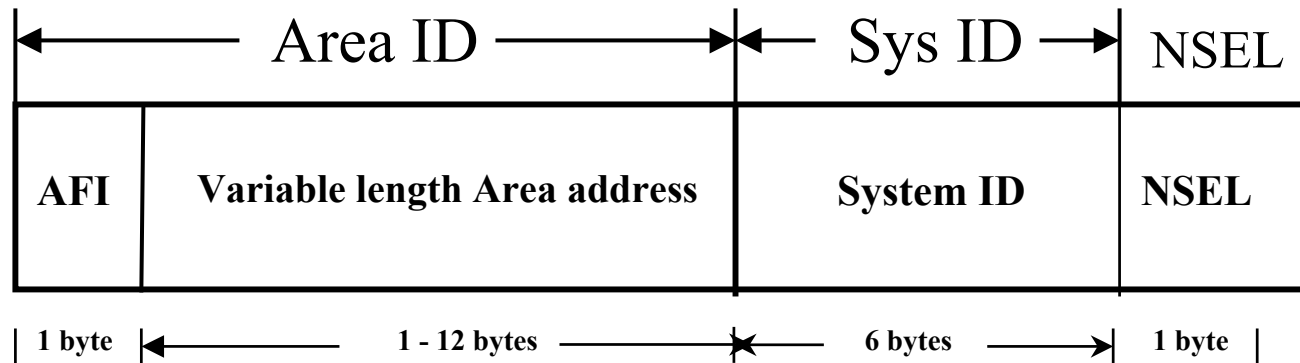
- ❑ The **Intermediate Systems to Intermediate System Routing Protocol (IS-IS)** was originally designed to route the ISO Connectionless Network Protocol (CLNP) . (ISO10589 or RFC 1142)
- ❑ Adapted for routing IP in addition to CLNP (RFC1195) as Integrated or **Dual IS-IS (1990)**
- ❑ IS-IS is a Link State Protocol similar to the Open Shortest Path First (OSPF). OSPF supports only IP
- ❑ IS-IS competed neck-to-neck with OSPF.
 - ❑ OSPF deployed in large enterprise networks
 - ❑ IS-IS deployed in several large ISPs

IS-IS Overview

- ❑ 3 network layer protocols play together to deliver the ISO defined Connectionless Network Service
 - ❑ CLNP
 - ❑ IS-IS
 - ❑ ES- IS - End System to Intermediate System Protocol
- ❑ All 3 protocols independently go over layer 2

CLNS Addressing

NSAP Format



- NSAP format has 3 main components
 - Area ID
 - System ID
 - N-Selector (NSEL) - value is 0x00 on a router
- NSAP of a router is also called a NET

CLNS Addressing

Requirements and Caveats

- ❑ At least one NSAP is required per node
- ❑ All routers in the same area must have a common Area ID
- ❑ Each node in an area must have a unique System ID
- ❑ All level 2 routers in a domain must have unique System IDs relative to each other
- ❑ All systems belonging to a given domain must have System IDs of the same length in their NSAP addresses

IS-IS Terminology

Intermediate system (IS) - *Router*

Designated Intermediate System (DIS) - *Designated Router*

Pseudonode - *Broadcast link emulated as virtual node by DIS*

End System (ES) - *Network Host or workstation*

Network Service Access Point (NSAP) - *Network Layer Address*

Subnetwork Point of attachment (SNPA) - *Datalink interface*

Packet data Unit (PDU) - *Analogous to IP Packet*

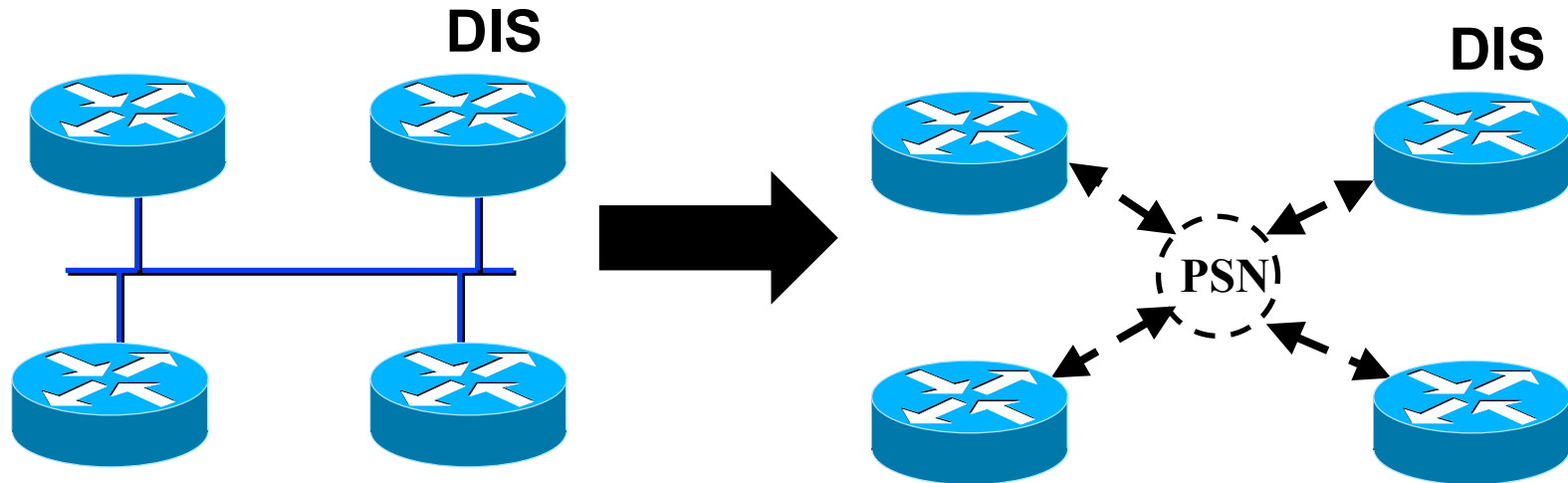
Link State PDU (LSP) - *Routing information packet*

Level 1 and Level 2 – *Area 0 and lower areas*

IS-IS Protocol Concepts: Network Nodes

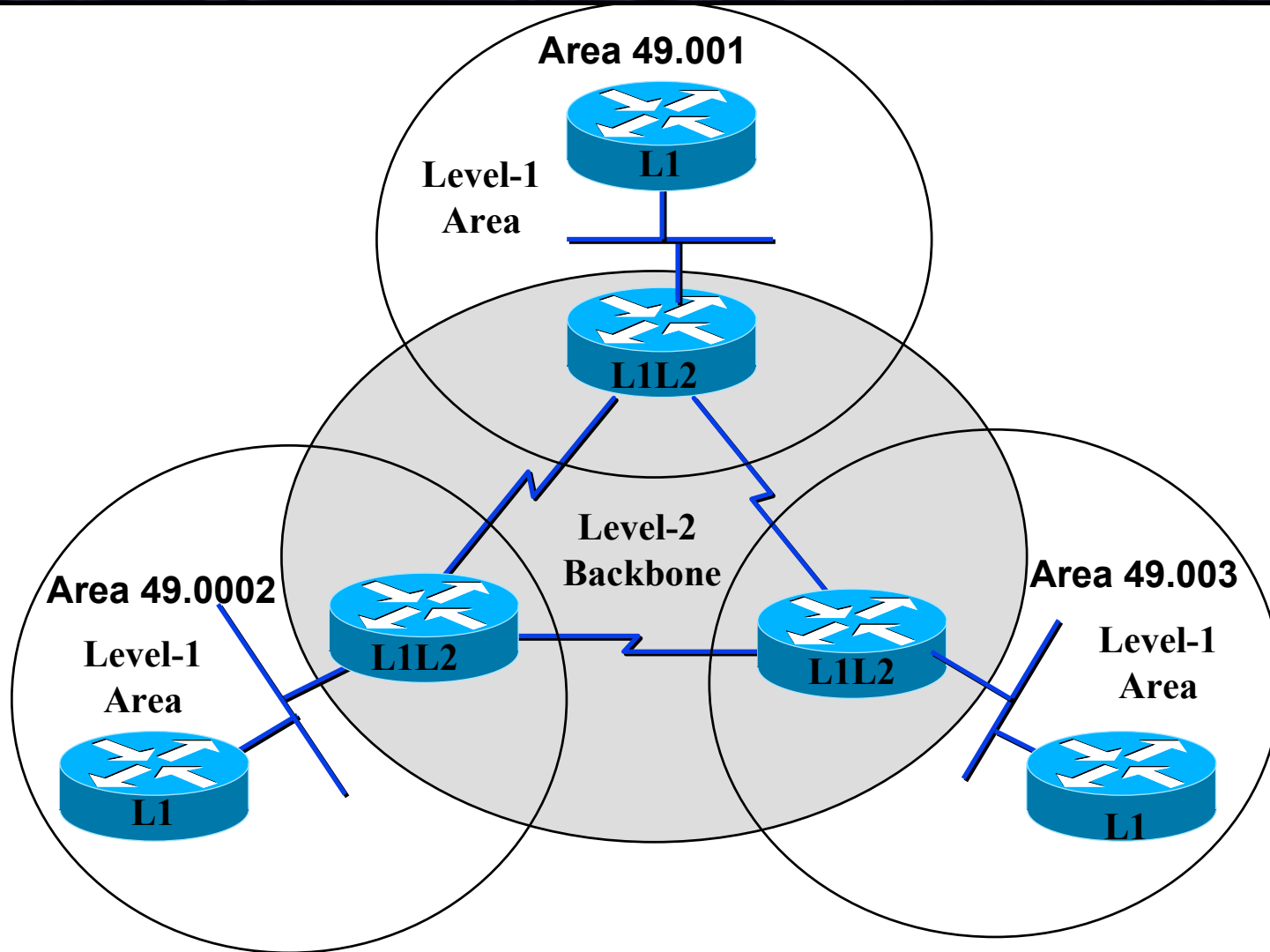
- ❑ Hosts
- ❑ Level-1 Routers
- ❑ Level-2 Routers
- ❑ Level-1 and Level-2 Pseudonodes on broadcast links only

IS-IS Protocol Concepts: Network Nodes

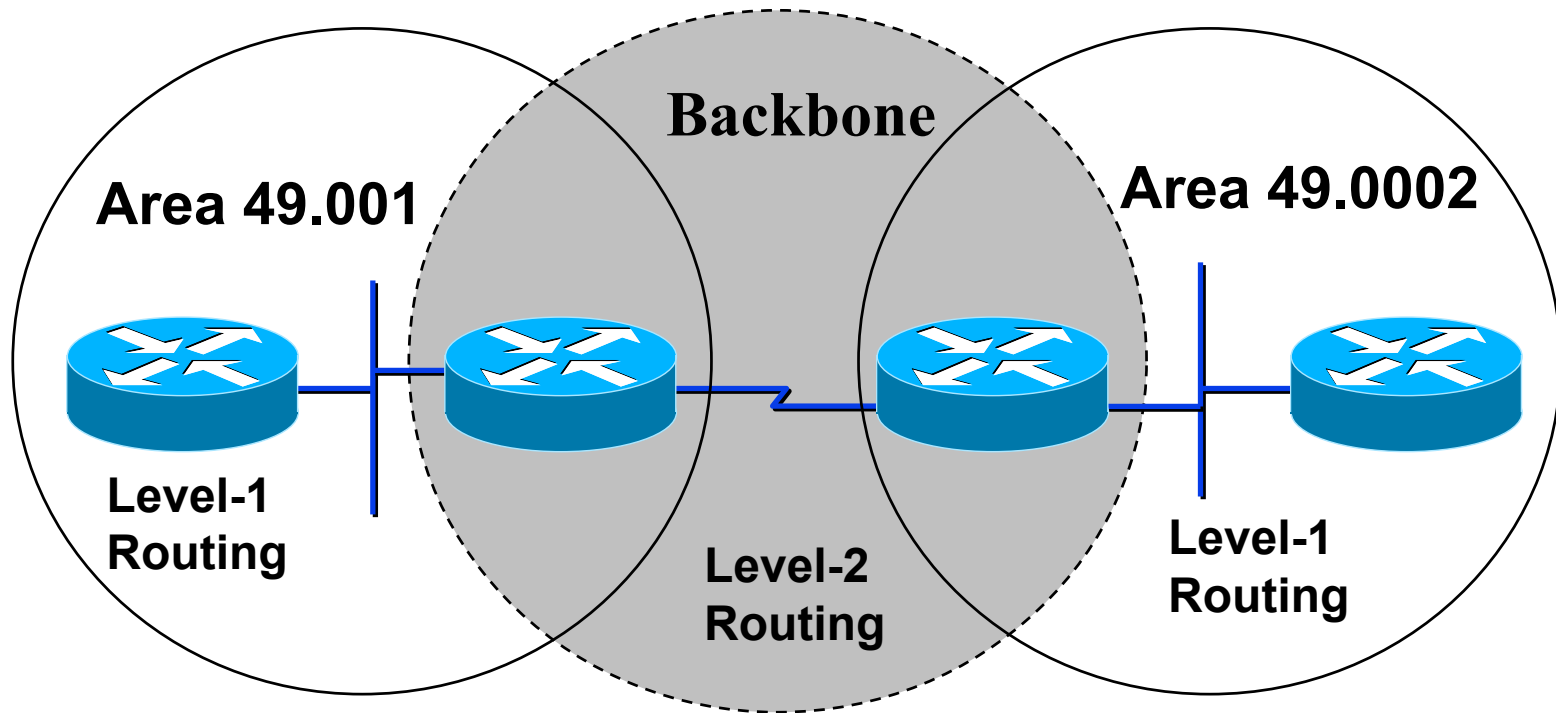


- ❑ Broadcast link represented as virtual node, referred to as **Pseudonode (PSN)**
- ❑ PSN role played by the **Designated Router (DIS)**
- ❑ DIS election is preemptive, based on interface priority with highest MAC address being tie breaker
- ❑ IS-IS has only one DIS. DIS/PSN functionality supports database synchronization between routers on a broadcast type link

IS-IS Protocol Concepts: Areas



IS-IS Protocol Concepts: Hierarchical Routing



- ❑ IS-IS supports 2-level routing hierarchy
- ❑ Routing domain is carved into areas. Routing in an area is level-1. Routing between areas is level-2
- ❑ All ISO 10589/RFC1195 areas are stubs

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IS-IS Protocol Concepts: IS-IS Packet Types

- ❑ **IS-IS Hello Packets (IIH)**
 - ❑ Level 1 LAN IS-IS Hello
 - ❑ Level 2 LAN IS-IS Hello
 - ❑ Point-to-point Hello
- ❑ **Link State Packets (LSP)**
 - ❑ Level 1 and Level 2
- ❑ **Complete Sequence Number packets (CSNP)**
 - ❑ Level 1 and Level 2
- ❑ **Partial Sequence Number Packets (PSNP)**
 - ❑ Level 1 and Level 2

IS-IS LS Database: IS-IS Packet Format

□ A Fixed Header

- Contains generic packet information and other specific information about the packet

□ Type, Length, Value (TLV) Fields

- TLVs are blocks of specific routing-related information in IS-IS packets

IS-IS LS Database: Generic Packet Format

				No. of Octets
Intra-domain Routing Protocol Discriminator				1
Length Indicator				1
Version/Protocol ID Extension				1
ID Length				1
R	R	R	PDU Type	1
Version				1
Reserved				1
Maximum Area Addresses				1
Packet-Specific Header Fields				
TLV Fields				

IS-IS LS Database: LSP Format

Intradomain Routing Proto Discriminator				Octets
Lenth Indicator				1
Version/Protocol ID Extension				1
ID Length				1
R	R	R	PDU Type	1
Version				1
Reserved				1
Maximum Area Addresses				2
PDU Length				2
Remaining Lifetime				ID Length + 2
LSP ID				4
Sequence Number				2
Checksum				1
P	ATT	LSPDBOL	IS Type	Variable
TYPE LENGTH VALUE FIELDS				

Level-1 TLVs

TLV Name	Type	Origin
Area Address	1	ISO 10589
Intermediate System Neighbors	2	ISO 10589
End System Neighbors	3	ISO 10589
Authentication information	10	ISO 10589
IP Internal Reachability Information	128	RFC 1195
Protocols Supported	129	RFC 1195
IP Interface Address	132	RFC 1195

Level-2 TLVs

TLV Name	Type	Origin
Area Address	1	ISO 10589
Intermediate System Neighbors	2	ISO 10589
Partition Designated Level-2 IS	4	ISO 10589
Prefix Neighbors	5	ISO 10589
Authentication information	10	ISO 10589
IP Internal Reachability Information	128	RFC 1195
Protocols Supported	129	RFC 1195
IP External Reachability Information	130	RFC 1195
Inter-domain Routing Protocol Information	131	RFC 1195
IP Interface Address	132	RFC 1195

High-level Comparison w/ OSPF

- ❑ Protocols are recognizably similar in function and mechanism (common heritage)
- ❑ Link state algorithms
- ❑ Two level hierarchies
- ❑ Designated Router on LANs
- ❑ Widely deployed (ISPs vs enterprises)
- ❑ Multiple interoperable implementations
- ❑ OSPF more “optimized” by design (and therefore significantly more complex)
- ❑ IS-IS not designed from the start as an IP routing protocol (and is therefore a bit clunky in places)

Detailed comparison points

❑ Encapsulation

- ❑ OSPF runs on top of IP=> Relies on IP fragmentation for large LSAs
- ❑ IS-IS runs directly over L2 (next to IP) => fragmentation done by IS-IS

❑ Media support

- ❑ Both protocols support LANs and point-to-point links in similar ways
- ❑ IS-IS supports NBMA in a manner similar to OSPF pt-mpt model: as a set of point-to-point links
- ❑ OSPF NBMA mode is configuration-heavy and risky (all routers must be able to reach DR; bad news if VC fails)

Comparison: Packet Encoding

- ❑ OSPF is “efficiently” encoded
 - ❑ Positional fields, 32-bit alignment
 - ❑ Only LSAs are extensible (not Hellos, etc.)
 - ❑ Unrecognized types not flooded. Opaque-LSAs recently introduced.
- ❑ IS-IS is mostly Type-Length-Value (TLV) encoded
 - ❑ No particular alignment
 - ❑ Extensible from the start (unknown types ignored but still flooded)
 - ❑ All packet types are extensible
 - ❑ Nested TLVs provide structure for more granular extension

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Comparison: Area Architecture

- ❑ Both protocols support two-level hierarchy of areas
- ❑ OSPF area boundaries fall within a router
 - ❑ Interfaces bound to areas
 - ❑ Router may be in many areas
 - ❑ Router must calculate SPF per area
- ❑ IS-IS area boundaries fall on links
 - ❑ Router is in only one area, plus perhaps the L2 backbone (area)
 - ❑ Biased toward large areas, area migration
 - ❑ Little or no multilevel deployment (large flat areas work so far)

Comparison: Database Granularity

- ❑ OSPF database node is an LSAdvertisement
 - ❑ LSAs are mostly numerous and small (one external per LSA, one summary per LSA)
 - ❑ Network and Router LSAs can become large
 - ❑ LSAs grouped into LSUUpdates during flooding
 - ❑ LSUUpdates are built individually at each hop
 - ❑ Small changes can yield small packets (but Router, Network LSAs can be large)

Comparison: Database Granularity

- ❑ IS-IS database node is an LSPacket
 - ❑ LSPs are clumps of topology information organized by the originating router
 - ❑ Always flooded intact, unchanged across all flooding hops (so LSP MTU is an architectural constant--it must fit across all links)
 - ❑ Small topology changes always yield entire LSPs (though packet size turns out to be much less of an issue than packet count)
 - ❑ Implementations can attempt clever packing

Comparison: Neighbor Establishment

- ❑ Both protocols use periodic multicast Hello packets, “I heard you” mechanism to establish 2-way communication
- ❑ Both protocols have settable hello/holding timers to allow tradeoff between stability, overhead, and responsiveness
- ❑ OSPF requires hello and holding timers to match on all routers on the same subnet (side effect of DR election algorithm) making it difficult to change timers without disruption
- ❑ IS-IS requires padding of Hello packets to full MTU size under some conditions (deprecated in practice)
- ❑ OSPF requires routers to have matching MTUs in order to become adjacent (or LSA flooding may fail, since LSUpdates are built at each hop and may be MTU-sized)

Neighbor Adjacency Establishment

- ❑ OSPF uses complex, multistate process to synchronize databases between neighbors
 - ❑ Intended to minimize transient routing problems by ensuring that a newborn router has nearly complete routing information before it begins carrying traffic
 - ❑ Accounts for a significant portion of OSPF's implementation complexity
 - ❑ Partially a side effect of granular database (requires many DBD packets)
- ❑ IS-IS uses its regular flooding techniques to synchronize neighbors
 - ❑ Coarse DB granularity => easy (a few CSNPs)

Designated Routers and Adjacency

- ❑ Both protocols elect a DR on multi-access networks to remove $O(N^2)$ link problem and to reduce flooding traffic
- ❑ OSPF elects both a DR and a Backup DR, each of which becomes adjacent with all other routers
 - ❑ BDR takes over if DR fails
 - ❑ DRship is sticky, not deterministic
- ❑ In IS-IS all routers are adjacent (adjacency less stateful)
 - ❑ If DR dies, new DR must be elected, with short connectivity loss (synchronization is fast)
 - ❑ DRship is deterministic (highest priority, highest MAC address always wins)
 - ❑ DRship can be made sticky by cool priority hack (DR increases its DR priority)

Comparison: LAN Flooding

- ❑ OSPF uses multicast send, unicast ack from DR
 - ❑ Reduces flood traffic by 50% (uninteresting)
 - ❑ Requires per-neighbor state (for retransmissions)
 - ❑ Interesting (but complex) acknowledgement suppression
 - ❑ Flood traffic grows as $O(N)$
- ❑ IS-IS uses multicast LSP from all routers, CSNP from DR
 - ❑ Periodic CSNPs ensure databases are synced (tractable because of coarse database granularity)
 - ❑ Flood traffic constant regardless of number of neighbors on LAN
 - ❑ But big LANs are uninteresting

Comparison: Routes and Metrics

- ❑ IS-IS base spec used 6-bit metrics on links
 - ❑ Allowed an uninteresting SPF optimization (CPUs are fast these days)
 - ❑ Proved difficult to assign meaningful metrics in large networks
 - ❑ Wide metric extension fixes this
- ❑ Dual IS-IS spec advertises only default into L1 areas
 - ❑ Inter-area traffic routed sub-optimally
 - ❑ Route leaking extension addresses this

Comparison: Pragmatic Considerations

- ❑ OSPF is much more widely understood
 - ❑ Broadly deployed in enterprise market
 - ❑ Many books of varying quality available
 - ❑ Preserves our investment in terminology
- ❑ IS-IS is well understood within a niche
 - ❑ Broadly deployed within the large ISP market
 - ❑ Folks who build very large, very visible networks are comfortable with it