



QUERY BY TABLE

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Abstract

Querying any information system requires the knowledge of some formal language, making it inaccessible to computer-naïve potential users. We propose a new intuitive querying mechanism where the query is a (well-formed) table. We extract the underlying logical structure of the table to retrieve values from a database. Query tables are interpreted to perform simple SELECT & JOIN operations. We demonstrate that query tables with different layouts but with the same underlying logical structure yield correct answers. This approach can be extended to form complicated conditional queries and queries involving aggregates.

1. Introduction

Query Languages are high level programming languages that allow users to make queries into databases. Examples include SMARTS, the cheminformatics standard for a substructure search, XQuery, a query language for XML data sources and SQL for relational databases. The semantics of the query are defined by a formal syntax specific to that language. In contrast, Query By Example and Query By Browsing have a simple graphical user interface. We present here, Query by Table (QBT), an intuitive mechanism based on the idea that well-formed tables can represent queries to a relational database.

One interesting aspect of tables is that answers to certain kinds of queries seem most naturally expressed in tabular form. Consider, for example, the query: “How does the volume of U.S. exports to Mexico compare to those to Canada for the years 2002 and 2003?” This question can be formulated as a query in the form of the table.

Table I. Query Q1

Year	U.S. Exports to Canada	U.S. Exports to Mexico
2002	<to be filled>	<to be filled>
2003	<to be filled>	<to be filled>

QBT takes inputs in the form of a table shown above and fills in the values. The method is motivated by our ongoing research in Project TANGO (Table ANalysis for Generating Ontologies) – a collaborative project between Brigham Young University and Rensselaer Polytechnic Institute - where we have applied conceptual model-based data extraction and table recognition to understand a table’s structure and its conceptual content for semi-automatic ontology generation [1]. We have used ideas drawn from the project to develop a simple paradigm for information retrieval.

To use tables as queries to a database, we distinguish between a table’s physical structure and its logical structure. Users may formulate a query table’s layout in many ways yet request the same values. Our starting point is the formalism put forth by Wang in [2] to interpret the table, i.e., to recover category labels and the functions that map them to a set of domains. Table interpretation is not trivial. Experiments have shown that even human “experts” often disagree on the label-value pairs of a table [3]. The Wang notation for a query table is currently extracted by an interactive tool that requires very little user intervention for simple tables [4] and can be eventually automated.

The paper is organized as follows. We first describe the Wang notation for a table and derive it for an example table. We discuss the dimensional model of a database and its similarity to Wang Notation. The distinctions between query tables and normal tables are sketched and the QBT process is demonstrated. We conclude with examples of the different types of queries that the system has successfully answered,

followed by a brief discussion of queries that QBT cannot handle yet.

2. Wang Notation

Wang defines an *abstract table* as an abstract data type and its layout structure as the *presentation form* of a table. The logical structure consists of entries and labels. The organization of the labels is called a *frame*, and the number of categories in the frame is the *dimension* of the abstract table. Applying a layout specification to an abstract table generates a *concrete table*. Tabular abstraction separates the logical structure of a table from its layout structure. The advantage of tabular abstraction is that the tables can be manipulated independently of their layout structure and we can easily alter the layout of a table by associating different topologies with the logical structure. Informally, the Wang Notation consists of two components (C, δ) where C is a finite set of labeled domains and δ is a mapping from the tree paths labels (or headers) to the possible values.

Table II is a 3-dimensional table. The categories are Year, Term and Mark. The Wang category notation for the table is

$$(Year, \{(1991, \varphi), (1992, \varphi)\})$$

$$(Term, \{(Winter, \varphi), (Spring, \varphi), (Fall, \varphi)\})$$

$$(Mark, \{(Assignments, \{(Ass1, \varphi), (Ass2, \varphi), (Ass3, \varphi)\}), (Examinations, \{(Midterm, \varphi), (Final, \varphi)\})\})$$

Year is the first category with 1991 and 1992 as subcategories. *Term* is the next category with *Winter*, *Spring* and *Fall* as the subcategories. *Mark* is the last category with two subcategories (*Assignments* and *Examinations*). *Assignments* and *Examinations* have their own subcategories. The delta notation shows which category cells are related to each of the individual values within the table. The δ notation for the first row of the table is

$$\delta(\{Year.1991, Term.Winter, Mark.Assignments.Ass1\})=85$$

$$\delta(\{Year.1991, Term.Winter, Mark.Assignments.Ass2\})=80$$

$$\delta(\{Year.1991, Term.Winter, Mark.Assignments.Ass3\})=75$$

$$\delta(\{Year.1991, Term.Winter, Mark.Examinations.Midterm\})=60$$

$$\delta(\{Year.1991, Term.Winter, Mark.Examinations.Final\})=75$$

Some tables lack spanning labels. For example, Table II would still be considered “valid” even if the header “Mark” were absent from the table. But the logical structure requires a root for the column header tree paths. This requires the addition of what Wang

Table II. Wang table

Year	Term	Mark				
		Assignments			Examinations	
		Ass1	Ass2	Ass3	Midterm	Final
1991	Winter	85	80	75	60	75
	Spring	80	65	75	60	70
	Fall	80	85	75	55	80
1992	Winter	85	80	70	70	75
	Spring	80	80	70	70	75
	Fall	75	70	65	60	80

called “virtual header”. Virtual headers are implicit headers added by the user in order to complete the logical structure of the table to obtain its abstract notation. The absence of explicit headers is the major difference between high-level and low-level table interpretation.

Lopresti et al. in [5] discuss the similarity of the relational paradigm to Wang’s idea of abstract tables. A database table is a two-dimensional table with *attributes* in one dimension and *tuples* in the other. The Wang notation captures the idea of attributes in its C notation. The tuples are specified in the δ notation. To generate the Wang Notation for query tables, we use the Interactive Wang Notation Tool (WNT) developed by Piyushee Jha.

3. Dimensional Database

Query By Table retrieves values from a relational database, which is modeled dimensionally. The following paragraphs explain the concepts of Dimensional Data Modeling.

A dimensional database is like a database cube of n dimensions. Users can access a slice of the database along any of its dimensions. The dimensional model is also called the *star-join* schema. The central table is the only table in the schema with multiple joins connecting it to all the other tables. The central table is called the *fact table* and the other tables are called *dimension tables*. The dimension tables have a single join that attaches them to the fact table. The fact table stores the measures or facts, which represent quantitative or factual data about the subject. The measures are generally numeric and correspond to the *how much* or *how many* aspects of a question. A dimension represents a single set of objects or events in the real world. Each

dimension we identify for the data model is implemented as a dimension table. Dimensions are the qualifiers that give meaning to the fact table measures. They answer the *what*, *when*, and *where* aspects of a question. A dimensional approach simplifies access to data that we want to summarize or compare. We chose tables from the Canada Statistics Website to construct a simple dimensional database. The dimensions of this database were Region, Year and Gender. The facts were identified as Population, University Degrees, Employment, Incidence of Asthma, Number of Farm Operators, Farm Operator Median Age, Number of Divorced, Single and Widowed citizens... (20 in total).

4. Query By Table

A query table is a well-formed user-constructed table that encapsulates all the information internal to its structure and uses exact attribute names for the facts to be retrieved. A query table does not use any headings, captions or titles to convey information. The Query by Table system is currently implemented in MATLAB. It accepts an MS Excel query table as an input and retrieves the values requested by the user. The process has five steps:

- Derive the Wang Notation for the query table.
- Parse the Wang Notation of the input table.
- Identify the facts and dimensions of the query.
- Form SQL queries from the facts and dimensions.
- Plug the results back into the query table.

Consider the query shown in the 3-dimensional Table III. An Excel macro program converts it into the ASCII format required by the Wang Notation Tool. The user interacts with the tool to produce the Wang Notation shown below:

C Notation:

```
(Year, {(2001,phi),(2006,phi)})
(Gender, {(Male,phi),(Female,phi)})
(Region, {(Alberta, phi),(Manitoba, phi)})
(Statistics, {(Number_of_Farm_Operators,phi),
(Farm_Operator_Median_Age,phi)})
```

Delta Notation:

```
delta({Statistics.Number_of_Farm_Operators,
Year.2001, Gender.Male,Region.Alberta})=XX
delta({Statistics.Farm_Operator_Median_Age ,
Year.2001, Gender.Male, Region.Alberta})=XX
```

The Wang Category Notation of the query table is parsed using regular expressions. Dimensions are

stored as separate tables in the database, so we can query the database to determine if a header is actually a

Table III. Query Q2

Year	Gender	Region	Statistics	
			Number_of_ Farm_Operators	Farm_Operator_ Median_Age
2001	Male	Alberta	XX	XX
		Manitoba	XX	XX
	Female	Alberta	XX	XX
		Manitoba	XX	XX
2006	Male	Alberta	XX	XX
		Manitoba	XX	XX
	Female	Alberta	XX	XX
		Manitoba	XX	XX

dimension for the data value, if it is a fact, or if it is merely a dummy header. We initially consider four values –Year, Gender, Region and Statistics. On querying the database for Year, Gender and Region we conclude that they are keys. For each row in the delta notation, all the subcategories associated with the keys are stored separately. Querying the fact table for Statistics retrieves no values. We conclude that the term is a dummy header and examine its subcategories. The database is queried for the first subcategory of Statistics, Number_of_Farm_Operators. Since it is a column in the fact table, it must be a fact and a value to be retrieved. Similarly, another fact is identified to be Farm_Operator_Median_Age for each combination of Region, Year and Gender. Now, we have the facts to be retrieved and the dimensions for the facts. This information is sufficient to form SQL queries to retrieve values from the database by looping through simple SELECT statements to perform a simple JOIN. The data is plugged back into the Excel sheet by making use of its explicit table layout.

4.1. Why Wang Notation for QBT?

Table IV. Query 3

Region	Gender	Number_Of_ Farm_Operators		Farm_Operator_ Median_Age	
		2001	2006	2001	2006
Alberta	Male	XX	XX	XX	XX
	Female	XX	XX	XX	XX
Manitoba	Male	XX	XX	XX	XX
	Female	XX	XX	XX	XX

Table IV is a variation of Table III, with one of the dimensions hidden under the facts to be retrieved. This table requires two additional virtual headers: Year, which is a suitable header name for the sub-categories 2001 and 2006 and Statistics, which is a reasonable header name for the subcategories Number

_of_Farm_Operators and Farm_Operator_Median_Age. Thus, after the addition of the Virtual Headers the Wang Notation becomes

```
(Region, {(Alberta,phi),(Manitoba,phi)})
(Gender, {(Male,phi),(Female,phi)})
(Year, {(2001,phi),(2006,phi)})
(Statistics, {(Number_of_Farm_Operators,phi),
(Farm_Operator_Median_Age,phi)})
delta({Statistics.Number_of_Farm_Operators,
Year.2001, Gender.Male, Region.Alberta})=XX
delta({Statistics.Farm_Operator_Median_Age,
Year.2001, Gender.Male, Region.Alberta})=XX
...
```

Although the layout of Table IV is different from the layout of Table III, the Wang Notation preserves the underlying logical structure. Even if the added virtual headers are different from Year and Statistics (say Period and Values), we can still retrieve the values by looking into their subcategories.

4.2. Example Queries

This section illustrates some of the queries that the system has successfully answered. We can present here only small examples. Some of our experimental query tables are much larger, but have similar structures.

Query Q4:

“Did Quebec have a larger fraction of single men & women than Manitoba in the years 2005 and 2007?”

Table V. Query Q4

Year	Region	Population		Marital_Status_Single	
		Male	Female	Male	Female
2005	Quebec	XX	XX	XX	XX
	Manitoba	XX	XX	XX	XX
2007	Quebec	XX	XX	XX	XX
	Manitoba	XX	XX	XX	XX

Query Q5:

“How did the Incidence of Diabetes and Asthma in Ontario & Nova Scotia vary in the years 2003 and 2005 among men and women?”

Table VI. Query Q5

Disease	Ontario				Nova Scotia			
	Male		Female		Male		Female	
	2003	2005	2003	2005	2003	2005	2003	2005
Incidence_of_Asthma	XX	XX	XX	XX	XX	XX	XX	XX
Incidence_of_Diabetes	XX	XX	XX	XX	XX	XX	XX	XX

Query Q6:

“How did university education affect employment for men and women in Ontario, Nunavut and Nova Scotia in 2006 and 2007?”

Table VII. Query Q6

Year	Region	Male		Female	
		University_Degrees	Employed	University_Degrees	Employed
2006	Ontario	XX	XX	XX	XX
	Nunavut	XX	XX	XX	XX
	Nova Scotia	XX	XX	XX	XX
2007	Ontario	XX	XX	XX	XX
	Nunavut	XX	XX	XX	XX
	Nova Scotia	XX	XX	XX	XX

5. Summary and Conclusions

We have demonstrated the retrieval of values from a simple dimensional database by means of a query table. The database is constructed by agglomerating information from web tables, but a query may require data from several tables. We are currently scaling up the database with many more tables. The challenge is to identify the dimensions and facts correctly and efficiently. We are now developing methods to answer more complicated queries including aggregation operations like summing, averaging and evaluating maxima, minima and conditional queries. The mild constraints imposed on the query tables are readily accepted by naïve users. Through our ongoing research in the TANGO project, we also expect to make query tables even more intuitive and more widely accessible.

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